



The following slides contain advanced material and are optional.

Outline



- Invariants
- Violating the invariant
- Marriage problem

Invariants explained in 60 seconds



- Consistency requirements for a class
- Established after object creation
- Hold, when an object is *visible*
 - Entry of a routine
 - Exit of a routine

```
class
  ACCOUNT
feature
  balance: INTEGER
invariant
  balance >= 0
end
```

Temporary violation



- Invariants can be violated temporarily
e.g. on object creation
- In Eiffel, invariants are checked on entry and exit of
a **qualified feature call**
- One exception: for calls to creation procedures,
invariants are not checked on entry to routine
 - e.g. *create cell.set_item (1)*
 - But checked for normal call: *cell.set_item (1)*
- See demo.

Public interface of person (without contracts)



```
class
  PERSON
feature
  spouse: PERSON
    -- Spouse of Current.

  marry (a_other: PERSON)
    -- Marry `a_other`.
  do
    ...
  end
end
```

```
class
  MARRIAGE
feature
  make
    local
      alice: PERSON
      bob: PERSON
    do
      create alice
      create bob
      bob.marry (alice)
    end
  end
end
```

Write the contracts

Hands-On

class *PERSON*

feature

spouse: PERSON

marry (a_other: PERSON)

require

??

ensure

??

invariant

??

end

A possible solution



class *PERSON*

feature

spouse: PERSON

marry (a_other: PERSON)

require

a_other /= Void

a_other /= Current

a_other.spouse = Void

spouse = Void

ensure

spouse = a_other

a_other.spouse = Current

end

invariant

spouse /= Void implies spouse.spouse = Current

spouse /= Current

end

Implementing *marry*



```
class PERSON
feature
  spouse: PERSON

  marry (a_other: PERSON)
    require
      a_other /= Void
      a_other /= Current
      a_other.spouse = Void
      spouse = Void
    do
      ??
    ensure
      spouse = a_other
      a_other.spouse = Current
    end

invariant
  spouse /= Void implies spouse.spouse = Current
  spouse /= Current
end
```

Implementing *marry* I

Hands-On

```
class PERSON
feature
  spouse: PERSON

  marry (a_other: PERSON)
    require
      a_other /= Void
      a_other /= Current
      a_other.spouse = Void
      spouse = Void
    do
      a_other.spouse := Current
      spouse := a_other
    ensure
      spouse = a_other
      a_other.spouse = Current
    end

invariant
  spouse /= Void implies spouse.spouse = Current
  spouse /= Current
end
```

Compiler Error:

No assigner
command for
a_other

Implementing *marry* II

Hands-On

```
class PERSON
feature
  spouse: PERSON

  marry (a_other: PERSON)
    require
      a_other /= Void and a_other /= Current
      a_other.spouse = Void
      spouse = Void
    do
      a_other.set_spouse (Current)
      spouse := a_other
    ensure
      spouse = a_other
      a_other.spouse = Current
    end

    set_spouse (a_person: PERSON)
    do
      spouse := a_person
    end

invariant
  spouse /= Void implies spouse.spouse = Current
  spouse /= Current
end
```

```
local
  bob, alice: PERSON
do
  create bob; create alice
  bob.marry (alice)
  bob.set_spouse (Void)
  -- What about the
  invariants
  -- of bob and alice?
end
```

Implementing *marry* III

Hands-On

```
class PERSON
feature
  spouse: PERSON

  marry (a_other: PERSON)
    require
      a_other /= Void and a_other /= Current
      a_other.spouse = Void
      spouse = Void
    do
      a_other.set_spouse (Current)
      spouse := a_other
    ensure
      spouse = a_other
      a_other.spouse = Current
    end

feature {PERSON}
  set_spouse (a_person: PERSON)
    do
      spouse := a_person
    end

invariant
  spouse /= Void implies spouse.spouse = Current
  spouse /= Current
end
```

What about the invariant of *a_other* in feature *marry*?

Implementing *marry*: final version



```
class PERSON
feature
  spouse: PERSON

  marry (a_other: PERSON)
    require
      a_other /= Void
      a_other.spouse = Void
      spouse = Void
    do
      spouse := a_other
      a_other.set_spouse (Current)
    ensure
      spouse = a_other
      a_other.spouse = Current
    end

feature {PERSON}
  set_spouse (a_person: PERSON)
    do
      spouse := a_person
    end

invariant
  spouse /= Void implies spouse.spouse = Current
  spouse /= Current
end
```

Ending the marriage

Hands-On

```
class PERSON
```

```
feature
```

```
  spouse: PERSON
```

```
  divorce
```

```
    require
```

```
      spouse /= Void
```

```
    do
```

```
      spouse.set_spouse (Void)
```

```
      spouse := Void
```

```
    ensure
```

```
      spouse = Void
```

```
      (old spouse).spouse = Void
```

```
  end
```

```
invariant
```

```
  spouse /= Void implies spouse.spouse = Current
```

```
  spouse /= Current
```

```
end
```

Is the order of instructions in divorce important for the invariant?

What we have seen



- Invariant should only depend on Current object
- If invariant depends on other objects
 - Take care **who can change state**
 - Take care in **which order** you change state
- Invariant can be temporarily violated
 - You can still call **features on Current object**
 - Take care in calling **other objects, they might call back**

Although writing invariants is not that easy, they are necessary to do formal proofs. This is also the case for loop invariants (which will come later).