ConAir: Featherweight Concurrency Bug Recovery Via Single-Threaded Idempotent Execution

Wei Zhang Marc de Kruijf Ang Li Shan Lu Karthikeyan Sankaralingam

ASPLOS 2013

CCC 2014 Andreas Dillier

Overview

- Motivation
- ConAir System
- Evaluation
- Criticism

Motivation - Overview

- Concurrency bugs remain hidden
- Severe failure
- Hard to fix

Motivation – Automatic Tools

- Compatibility
- Correctness
- Generality
- Performance

Motivation – Automatic Fixing

- Additional Synchronisation
- Needs to know about the root cause

Motivation – Prohibit Interleaving

- Additional serialisation leads to performance loss
- Only works on certain interleavings
- Some approaches need programmer annotations

Motivation – Traditional Rollback Recovery

- Require checkpointing
- All threads are rolled back
- Require OS/hardware modification to run efficiently

ConAir - Overview

- Single thread rollback
- No checkpoints
- Static program modification

- Atomicity violations
 - About 70% of non-deadlock bugs
- Rollback of one thread establishes serialised execution
- About 92% recoverable

• Write after write (WAW)

Log = CLOSE; Log = OPEN; If (Log != OPEN) { //output failure}

• Write after write (WAW)

Log = CLOSE; Log = OPEN; If (Log != OPEN) { //output failure}

Rollback of this thread recovers the program

• Read after write (RAW)

ptr = aptr; tmp = *ptr; ptr = NULL;

Read after write (RAW)

ptr = NULL;

Rollback of this thread recovers the program

Read after read (RAR)

if (ptr)
{ fputs(ptr); }

ptr = NULL;

Read after read (RAR)

Rollback of this thread recovers

ptr = NULL;

• Write after read (WAR)

count += deposit1;
printf(cnt);

cnt += deposit2;

• Write after read (WAR)

```
count += deposit1;
printf(cnt);
```

cnt += deposit2;

Rollback of this thread recovers

Order-violation

- 30% of non-deadlock bugs
- About 50% recoverable

SomeClass b = NULL; b = new SomeClass(); b.foo();

Order-violation

- 30% of non-deadlock bugs
- About 50% recoverable

SomeClass b = NULL; b = new SomeClass(); b.foo();

Rollback of this thread recovers

- Deadlock bugs
 - About 40% of bugs
- Recovers if one thread rolled back

ConAir - Basics

- Identify potential failure sites
- Identify idempotent region for each failure site
- Insert recovery code
- Two modes: Fix and Survival

ConAir – Failure Site Identification

Survival mode

- Assertions
 - \rightarrow Can use assertion to indicate output failure
- Heap/global pointer dereference
- Deadlock detection with any detection tool

ConAir – Failure Site Identification

• Fix mode

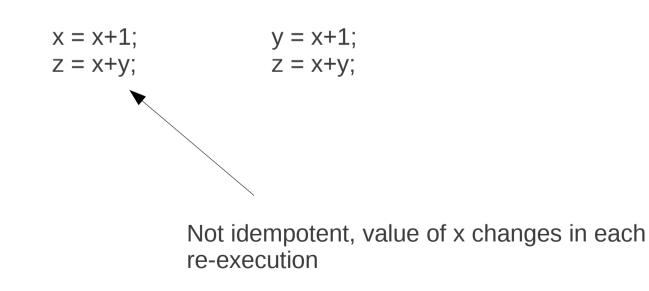
- Programmer specifies the location of the failure

- Re-execution only on idempotent regions
- Guarantees correctness
- May be too weak for many bugs
 - But has a low overhead

- No writes to shared Variables
- No I/O
- No idempotent destroying write to local variables

$$x = x+1;$$
 $y = x+1;$
 $z = x+y;$ $z = x+y;$

- No writes to shared Variables
- No I/O
- No idempotent destroying write to local variables



- Discovery non-trivial
 - Source code ↔ bit code ↔ binary code
- Binary code analysis alone complicated
- Search all backward paths from failure sites

- Weaken the definition
 - Idempotent function calls
 - Parent functions
- Requires more analysis

- At start of re-execution region: setjmp
 - Saves the register image
- At failure site: longjmp
 - Loads the register image and executes from setjmp
- Multiple retries

• Original code

... if(e){ ... } else {

}

__assert_fail(...);

Modified code

```
_thread jmp_buf c;
  _thread int RetryCnt=0;
Reexecution:
    setjmp(c);
     ... //idempotent region
    if(e){
     . . .
    } else {
Failure:
         while(RetryCnt++<maxRetryNum){</pre>
              longjmp(c, 0);
            _assert_fail(...);
     }
```

Thread local variables for registers and retry count

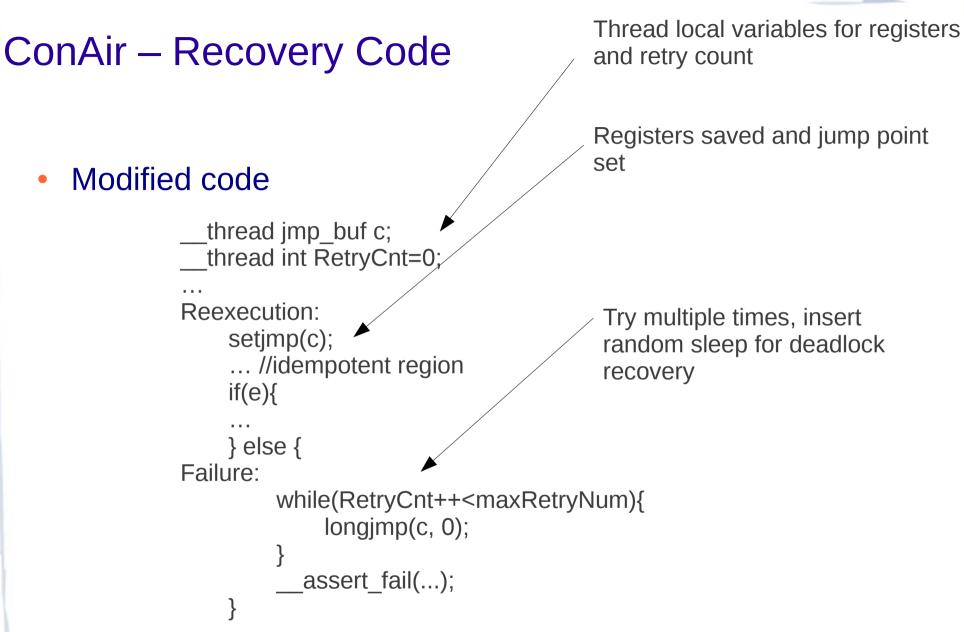
```
Modified code
             thread jmp_buf c;
             thread int RetryCnt=0;
          Reexecution:
               setjmp(c);
               ... //idempotent region
               if(e){
                . . .
               } else {
          Failure:
                    while(RetryCnt++<maxRetryNum){</pre>
                         longjmp(c, 0);
                       _assert_fail(...);
               }
```

Modified code

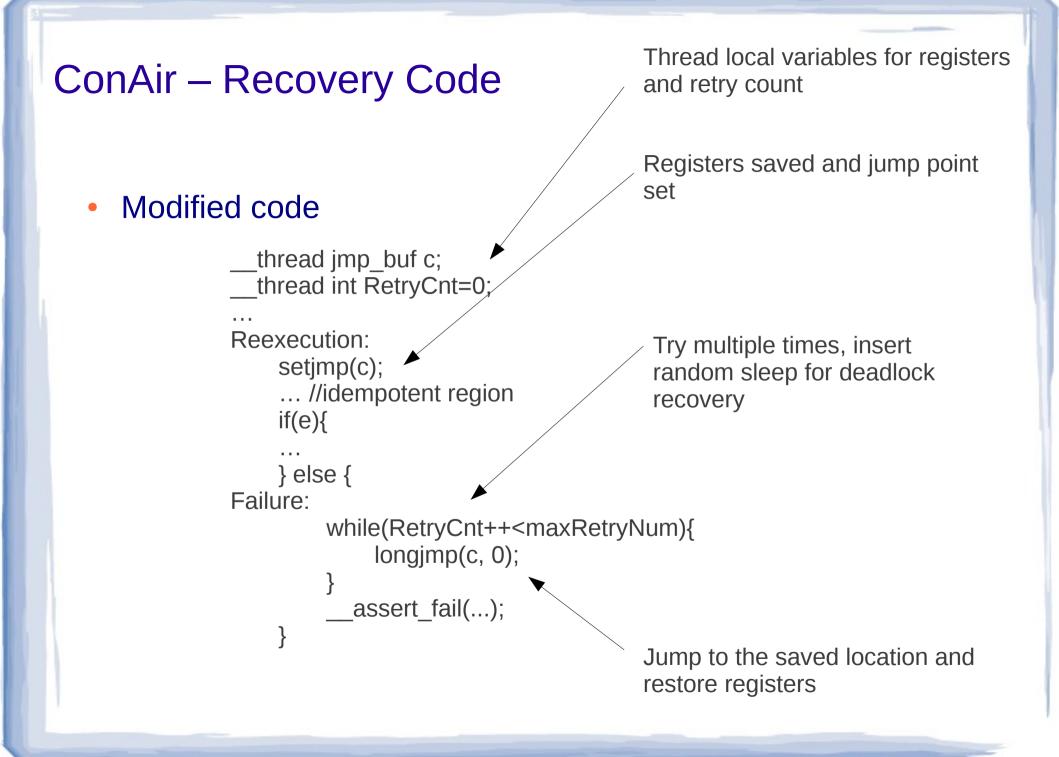
Thread local variables for registers and retry count

Registers saved and jump point set

```
thread jmp_buf c;
  _thread int RetryCnt=0;
Reexecution:
    setjmp(c);
     ... //idempotent region
    if(e){
     . . .
    } else {
Failure:
         while(RetryCnt++<maxRetryNum){</pre>
              longjmp(c, 0);
            assert_fail(...);
```



Try multiple times, insert random sleep for deadlock



ConAir - Optimisations

- Allow library functions
 - Extend idempotent region
 - Needed for certain recoveries (Deadlock)
 - Need compensation function (lock/unlock, malloc/free)
- ConAir allows malloc and lock in idempotent regions
 - Cleanup before longjmp

ConAir – Optimisations

• Remove code from unrecoverable fail sites

- Statically proven
- Deadlock recovery with no lock in idempotent region
- Non-deadlock recovery with no shared-variable reads

ConAir - Optimisations

- Include parent functions in idempotent region
- Should at least change one argument
- Significant overhead in static analysis

Evaluation

- 10 bugs in open-source libraries
- Wide variety of root causes and failure symptoms
- Analyze performance, overhead and recovery time
- Also analyze static analysis time

Evaluation

- Modify buggy code with sleep instructions
 - Almost 100% failure rate
- Run 1000 times with applied ConAir
- Successfully recovered if none causes the bug

Evaluation

• Run time overhead measured on the original source code

Results – Fix Mode

- No measurable overhead
 - Small number of failure sites
- Recovered all failures

Results – Survival Mode

- Small overhead (<1%)
- Could recover 8/10 bugs
 - I/O operations would require program annotations

Results – Recovery Time

- At most 17 milliseconds
- Much better than crash/program restart

Results – Static Analysis

- FFT (1.2K lines of code)
 - Less than a second
- MySQL(~685k lines of code)
 - 4 hours
- Inter procedural analysis requires big part
 - Only 50 seconds in MySQL are spent on intra procedural analysis

Limitations

- No completeness
 - No completeness

Criticism

- Aims of ConAir are met
- Surprisingly low overhead
- Easy to read, enough explanations
- Hard to find negative points, since ConAir is more of a heuristic approach