





Part 1: Language constructs

# **1.3 EXCEPTION HANDLING**

```
public class Printer {
   public print(int i) {
      try {
        throw new Exception()
      }
      catch(Exception e) { }
      }
```

# **Eiffel: Exception Handling**

class PRINTER feature print\_int (a\_int: INTEGER) local I retried: BOOLEAN do if not / retried then (create {DEVELOPER EXCEPTI else -- Do something end rescue *I* retried := **True** retry end end OLD VIOLATION

EXCEPTION 🖃 🎱 ASSERTION VIOLATION CHECK\_VIOLATION INVARIANT\_VIOLATION LOOP INVARIANT VIOLATION POSTCONDITION VIOLATION PRECONDITION VIOLATION VARIANT\_VIOLATION DEVELOPER EXCEPTION 🖃 🎱 MACHINE EXCEPTION 🖃 🎱 HARDWARE EXCEPTION FLOATING\_POINT\_FAILURE 🖃 🍉 OPERATING\_SYSTEM\_EXCEPTION COM\_FAILURE OPERATING\_SYSTEM\_FAILURE OPERATING SYSTEM SIGNAL FAILURE OBSOLETE\_EXCEPTION EXCEPTION\_IN\_SIGNAL\_HANDLER\_FAILURE RESCUE\_FAILURE RESUMPTION FAILURE 🖃 🎱 SYS\_EXCEPTION EIFFEL\_RUNTIME\_PANIC 🖃 🎱 EIF EXCEPTION 🖃 🎱 EIFFEL\_RUNTIME\_EXCEPTION 🖃 🍉 DATA\_EXCEPTION IO FAILURE MISMATCH FAILURE SERIALIZATION\_FAILURE EXTERNAL FAILURE NO\_MORE\_MEMORY 🖃 🎱 LANGUAGE EXCEPTION BAD INSPECT VALUE EIFFELSTUDIO\_SPECIFIC\_LANGUAGE\_EXCEPTION ADDRESS APPLIED TO MELTED FEATURE CREATE\_ON\_DEFERRED ROUTINE FAILURE VOID\_ASSIGNED\_TO\_EXPANDED VOID TARGET

```
feature
   transmit (a_p: PACKET)
        -- transmit packet a p
     local
        I current retries: INTEGER
        r: RANDOM_NUMBER_GENERATOR
      do
        line.send (a_p)
      rescue
        if I current retries < max retries then
           r.next
            wait_millisecs (r.value_between(20, 50))
            current retries := current retries + 1
           retry
        end
   end
end
```

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# **1.4 ONCE ROUTINES**

```
foo: INTEGER

once

Result := factorial (10)

end

test_foo

do

io.put_integer (foo) -- 3628800, calculated

io.put_integer (foo) -- 3628800, directly returned

end
```

- Executed the first time
- Result is stored
- In further calls, stored result is returned
- In other languages
  - Static variables
  - Singleton pattern



Constants, other than basic types *i: COMPLEX*once create Result.make (0, 1) end

Lazy initialization settings: SETTINGS once create Result.load\_from\_filesystem end

>Initialization procedures
 init\_graphics\_system
 once ... end

# Part 1: Language constructs





# **Style rule**

# For indentation, use tabs, not spaces





# More style rules

- Class name: all upper-case Full words, no abbreviations (with some exceptions)

- Classes have global namespace: two classes cannot have the same name (even in different clusters)

- Usually, classes are prefixed with a library prefix

EiffelVision2: EV\_

Base is not prefixed



class > PREVIEW inherit TOURISM feature explore -- Show city info -- and route. do Paris.display Louvre.spotlight Line8.highlight Route1.animate end end

# **Even more style rules**



# **Eiffel Naming: Locals / Arguments**

Locals and arguments share namespace with features

Name clashes arise when a feature is introduced, which has the same name as a local (even in parent)

#### ➤To prevent name clashes:

- Locals are prefixed with
- Some exceptions like "i" exist
- Arguments are prefixed with a\_



Part 1: Language constructs

# **1.6 GENERICS**

#### class MY QUEUE [G]

#### feature

#### item: G

-- First item in queue.

do ... end

extend (a\_element: G) -- Add new element. do ... end

end



G is called the generic parameter. By convention, the generic parameter name is G. If there are more parameters, use G, H, etc. or a meaningful abbreviation such as K for keys in a hash table

# **Creating instances of generics classes**

#### class EXAMPLE1

#### feature

int\_queue -- An integer queue. local qi: MY\_QUEUE [INTEGER] do create qi qi.extend (35) qi.extend (6) end end

class EXAMPLE2

#### feature string\_queue -- A string queue. local qs: MY\_QUEUE [STRING] do create qs qs.extend ("Asterix") qs.extend ("Obelix") qs.extend ("Suffix") end end

#### class

MY\_LIST [G -> COMPARABLE]

#### feature

#### item: G

-- First item in queue.

do ... end

. . .

```
extend (a_element: G)
-- Add new element.
do
```

```
... if a_element < item then
```

The generic parameter G must be a class inheriting from COMPARABLE



end

### **Creating instances of constraint generics classes**

. . .

-- Valid declarations li: MY\_LIST [INTEGER] ls: MY\_LIST [STRING] lr: MY\_LIST [REAL] ld: MY\_LIST [DOUBLE]

. . .

-- Invalid declarations la: MY\_LIST [ACCOUNT] lb: MY\_LIST [BANK] lm: MY\_LIST [MAIN]

> Classes ACCOUNT, BANK & MAIN don't inherit from COMPARABLE

Part 1: Language constructs

# **1.8 INFORMATION HIDING**



Procedure: doesn't return a result

- Yields a command
- Calls are instructions

Function: returns a result

- Yields a query
- Calls are expressions



# It doesn't matter to the client whether you look up or compute



### **Uniform Access: an example**

#### balance = list\_of\_deposits.total - list\_of\_withdrawals.total



A call such as your\_account.balance could use an attribute or a function

# Exporting (making public) an attribute

In Eiffel, exporting an attribute means exporting it read-only

From the outside, it is not shown as an attribute, just as a **query**: it could be a function

In C++, Java & C#, if you make public an attribute\* *x*, it is available for both read and write:

> a1.x := v

As a result, it is almost always a bad idea to export an attribute.

\* (field, member variable)

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## **Getter functions**

In C++, Java & C#, the standard technique, if *private\_x* is secret, is to export an associated **getter function**:

```
x : T
do
Result := private_x
end
```

Eiffel needs no getter functions: just export the attribute

This is safe: the attribute is exported

- Only for reading
- Without the information that it is an attribute: it could be a function (Uniform Access principle)

# **Information hiding**



#### Status of calls in a client with a1: A:

- a1.f, a1.g: valid in any client
- a1.h: invalid everywhere
  - (including in A's own text!)
- a1.j: valid only in B, C and their descendants (not valid in A!)
- a1.m: valid in B, C and their descendants, as well as in A and its descendants

Information hiding only applies to use by clients, i.e. using dot notation or infix notation, as with *a1.f* (*Qualified* calls).

Unqualified calls (within class) not subject to information hiding:

```
class A feature {NONE }
h do ... end
feature
```



# **PART 2: CONTRACTS**

# Contracts

A contract is a semantic condition characterizing usage properties of a class or a feature

Three principal kinds:

- Precondition
- Postcondition
- Class invariant

# **Design by Contract**

Together with the implementation ("*how*") of each software element, describe "*what*" it is supposed to do: its contract

Three basic questions about every software element:

What does it assume?

> What does it guarantee?
> What does it maintain?
Precondition
Postcondition

# **Contracts in programming languages**

Eiffel: integrated in the language

Java: Java Modeling Language (JML), iContract etc.

.Net languages: Code Contracts (a library)

Spec# (Microsoft Research extension of C#): integrated in the language

UML: Object Constraint Language

etc.

# Precondition

Property that a feature imposes on every client:



A client calling a feature must make sure that the precondition holds before the call

A client that calls a feature without satisfying its precondition is faulty (buggy) software.

#### Another example:

```
extend (a_element: G)

require

valid_elem: a_element /= void

not_full: not is_full

do ... end
```

```
A feature with a require clause

require

label_1: cond_1

label_2: cond_2 ...

label_n: cond_n

is equivalent to

require

label: cond_1 and cond_2 and ... cond_n
```





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### Let's code...

Go to:



https://codeboard.io/projects/86

*Task*: in class CUSTOMER, write a precondition for the creation routine make\_with\_name\_and\_age

*Task*: create an invalid CUSTOMER object and try to run your program. What happens?

*Task*: fix your **CUSTOMER** object to satisfy the precondition of the creation routine

# Postconditions

Precondition: obligation for clients Postcondition: benefit for clients

```
extend (a_element: G)
ensure
inserted: i_th (count) = a_element
```

```
index (a_element: G): INTEGER
```

#### ensure

```
exists: result > 0 implies i_th (result) = a_element
no_exists: result = -1 implies not is_inserted (a_element)
```

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### Let's code...

Go to:



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https://codeboard.io/projects/86

*Task*: in class CUSTOMER, write a postcondition for the creation routine make\_with\_name\_and\_age

*Task*: modify the implementation of make\_with\_name\_and\_age such that it breaks your postcondition. Run the program. What happens?

# **Old notation**

Usable in postconditions only

Denotes value of an expression as it was on routine entry

Example (in a class ACCOUNT):

```
balance: INTEGER
              -- Current balance.
deposit (v: INTEGER)
              -- Add v to account.
       require
              positive: v > 0
       do
       ensure
              added: balance = old balance + v
       end
```

A feature must make sure that, if its precondition held at the beginning of its execution, its postcondition will hold at the end.

A feature that fails to ensure its postcondition is buggy software.



## Invariant

An invariant states properties about an object that are true

- after the object has been initialized
- before and after every routine call (but not necessarily in between a call)

The invariant is listed after the last feature block.

A class with no **invariant** is the same a **invariant** always\_OK: *True* 

Example (from class ARRAY):

#### invariant

area\_exists: area /= Void consistent\_size: capacity = upper - lower + 1 non\_negative\_count: count >= 0 index\_set\_has\_same\_count: valid\_index\_set

#### class

BANK\_ACCOUNT create make feature make (n : STRING) -- Set up with name n

require n /= Void

do

name := n balance := 0

ensure name = n

end

name : STRING balance : INTEGER deposit (v: INTEGER) -- Add amount v do balance := balance + vensure balance = old balance + vend invariant name /= Void

*balance* >= 0

#### end

### Let's code...

Go to:



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https://codeboard.io/projects/86

*Task*: in class ACCOUNT, replace all -- Important: ... comments with contracts

# **Contracts and inheritance (Example)**

class	ACCOUNT
ACCOUNT_MANAGER	ACCOUNT_
feature Operations	MANAGER
<pre>init_new_account(a_acc: ACCOUNT)</pre>	
do all initialization of po	not strengthen ondition because lymorphism and mic binding.
	class
class ACCOUNT	SPECIAL_ACCOUNT inherit ACCOUNT redefine set_balance end
feature Operations	feature Operations
<pre>set_balance(a_balance: DOUBLE)     require     non_neg: a_balance &gt;= 0     do         balance := a_balance</pre>	<pre>set_balance(a_balance: DOUBLE)     require     min_bal: a_balance &gt; 100     do</pre>
end	<pre>balance := a_balance end</pre>

When redeclaring a routine, we may only:

Keep or weaken the precondition

Keep or strengthen the postcondition

Invariant Inheritance rule:

The invariant of a class automatically includes the invariant clauses from all its parents,

"and"-ed.

A simple language rule does the trick!

Redefined version may have nothing (assertions kept by default), or

require else new\_pre
ensure then new\_post

Resulting assertions are:

original\_precondition or new\_pre

original\_postcondition and new\_post

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