Mock Exam 1

ETH Zurich

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Name: 

Group: 

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1 Multiple choice (7.5 points)

Put checkmarks in the checkboxes corresponding to the correct statements. There is at least one correct answer per question. A correctly checked or unchecked box is worth 0.5 points. An incorrectly checked or unchecked box is worth 0 points. Completely unanswered questions are worth 0 points.

Example:

Which of the following statements are true?

a. The sun is a mass of incandescent gas. ☒ 0.5 points
b. $2 \times 4 = 8$ □ 0 points
c. “Rösti” is a kind of sausage. ☒ 0 points
d. C is an object-oriented programming language. □ 0.5 points

1. Control structures and recursion.
   a. If we know that a loop decreases its variant and that it never goes below 5, then we know that the loop terminates. □
   b. The loop invariant is checked at the end of loop initialization (before entering the loop itself). □
   c. The loop invariant tells us how many times the loop will be executed. □
   d. In Eiffel a procedure can have an empty body (do end). □
   e. A loop can always be rewritten as a finite sequence of conditional statements and compound statements. □

2. Objects and classes
   a. All entities store references to run-time objects. □
   b. Different entities can reference the same object. □
   c. Clients of a class $X$ can see all features declared in class $X$. □
   d. A class needs to tell its clients whether a query is an attribute or a function. □
   e. Objects can be created from every class. □

3. Design by Contract
   a. The creation procedure only needs to ensure that the invariant of the created object holds at the end of the procedure body. □
   b. Every procedure ensures that the postcondition True holds. □
   c. The class invariant needs to hold before every procedure call. □
   d. For functions, the precondition may not refer to the result expression and the postcondition may not refer to the arguments of the function. □
   e. A feature with precondition false is accepted by the compiler. □
2 Specifying Software through Contracts (14 points)

A range of integers can be conveniently represented using the boundary values of the range, e.g., the range of integers between \( m \) and \( n \) (inclusive) can be represented using \([m, n]\). Given a range \( R \), we use \( S_R \) to denote the set of integers within \( R \), i.e.

\[
S_{[m,n]} = \{ x \mid m \leq x \leq n \}.
\]

Listing 1 shows a class \( RANGE \), which abstracts integer ranges and provides functions that operate on them. The preconditions of the functions are already defined in the class; the function results, however, are only given in the comments in terms of the boundary values and the integer sets corresponding to the operand ranges. For example, the comment of function \( is\_equal \) stipulates that \( Result \) should be \( True \) if and only if \( Current \) and \( other \) represent the same set of integers, and the comment of function \( add \) specifies the integer set of \( Result \) should be equal to the union of the sets of \( Current \) and \( other \).

Read through the code, then complete the postconditions so that they reflect the function comments.

Please note:

- The number of dotted lines is not indicative of the number of missing contract clauses.
- You need to write \( True \) at places where you think no explicit contract is necessary: leaving a postcondition empty gives you 0 point for that section.
- The following features from class \( INTEGER \) may be useful:

```plaintext
class INTEGER

feature
  max (other: INTEGER): INTEGER
  min (other: INTEGER): INTEGER

end
```

Listing 1: Class \( RANGE \)

```plaintext
note
description: "A range of integers."

class RANGE

inherit

  ANY
  redefine is_equal end

create make

feature{NONE} -- Initialization
```
make \((l, r : \text{INTEGER})\)
\begin{verbatim}
  do
    left := l
    right := r
  end
\end{verbatim}

feature -- Access.

\textit{left} : \text{INTEGER}
  \begin{itemize}
  \item Lower boundary of the range.
  \item \(S_{\text{Current}} = \{x \mid \text{left} \leq x \leq \text{right}\}\)
  \end{itemize}

\textit{right} : \text{INTEGER}
  \begin{itemize}
  \item Upper boundary of the range.
  \item \(S_{\text{Current}} = \{x \mid \text{left} \leq x \leq \text{right}\}\)
  \end{itemize}

feature -- Query

\textit{is\_equal} (\textit{other}: like \text{Current}): \text{BOOLEAN}
  \begin{itemize}
  \item Result = \(S_{\text{Current}} = S_{\text{other}}\)
  \end{itemize}

require
  \textit{other} \neq \text{Void}

ensure

\textit{is\_empty}: \text{BOOLEAN}
  \begin{itemize}
  \item Result = \(S_{\text{Current}} = \emptyset\)
  \end{itemize}

require
  \text{True}

ensure

\textit{is\_sub\_range\_of} (\textit{other}: like \text{Current}): \text{BOOLEAN}
  \begin{itemize}
  \item Result = \(S_{\text{Current}} \subseteq S_{\text{other}}\)
  \end{itemize}

require
  \textit{other} \neq \text{Void}

ensure

\textit{is\_super\_range\_of} (\textit{other}: like \text{Current}): \text{BOOLEAN}
  \begin{itemize}
  \item Result = \(S_{\text{Current}} \supseteq S_{\text{other}}\)
  \end{itemize}

require
other /= Void
ensure


left_overlaps (other: like Current): BOOLEAN

−− Result = (left ∈ (S_{Current} ∩ S_{other}))
require
other /= Void
ensure


right_overlaps (other: like Current): BOOLEAN

−− Result = (right ∈ (S_{Current} ∩ S_{other}))
require
other /= Void
ensure


overlaps (other: like Current): BOOLEAN

−− Result = (S_{Current} ∩ S_{other} ≠ ∅)
require
other /= Void
ensure


feature −− Operation

add (other: like Current): RANGE

−− S_{Result} = (S_{Current} ∪ S_{other})
require
other /= Void
result_is_range : is_empty or other.is_empty or overlaps (other)
ensure
Result /= Void
subtract (other: like Current): RANGE

\[ S_{\text{Result}} = (S_{\text{Current}} - S_{\text{other}}) \]

require:
other /= Void
result_is_range : not overlaps (other)
or left_overlaps (other) or right_overlaps (other)

ensure
Result /= Void

end
3 Data Structures (15 points)

3.1 Background information

A skip list is a data structure that expands on the idea of a linked list. A node in a linked-list has 1 link; each node in a skip list has 4 links, up, down, left, and right.

A skip list has the following properties:

• The nodes are arranged into rows; each row is a list of sorted elements.

• Every row, except for the bottom row, contains a subset of the elements beneath it, as in Figure 1. This implies that the bottom row contains all the elements in the skip list.

• All nodes are mutually linked, i.e. node_a. left = node_b iff node_b. right = node_a, and likewise for up and down.

• Every row begins with a universal minimal element (represented here by $-\infty$).

• If an element exists in two adjacent rows, then the nodes are linked through the up/down attributes. This can be seen for the elements 20 and $-\infty$ in Figure 1.

![Figure 1: Initial skip list](image)

When a new element is inserted into the skip list, it is first inserted into the bottom row, as in Figure 2.

![Figure 2: Skip list after insertion of 30 into the bottom row](image)

Whenever a node is added to any row, there is a chance that it will be promoted, adding it to the row above, as in Figure 3. If there is no row above, a new one will be created. This promotion to the row above happens randomly, and a promotion can trigger another promotion (again, randomly).

3.2 Task

For the task the search feature is already implemented, and returns the rightmost node in the bottom row of the skip list less-than or equal to the argument elem. Feature is_promoted randomly returns True or False, indicating whether to promote a node at any given time. You must implement:
• *insert_in_row* (*a_pre*, *a_node*: *SKIP_LINKABLE) inserts *a_node* directly after *a_pre* with no promotion. An instance of this can be seen in the transformation between Figure 1 and Figure 2.

• *promote* (*a_link*: *SKIP_LINKABLE) takes *a_link*, which is already inserted in a row, and either promotes it or does nothing. Remember, *promote* can trigger another promotion.

• *insert* (*elem*: *INTEGER*) takes an element and inserts a new node into the correct position in the skip list, including promotion (if any).

While writing these procedures you are encouraged to use any applicable features already available in the *SKIP_LIST* and *SKIP_LINKABLE* classes (i.e. the features shown below without dotted lines).

```plaintext
class
  SKIP_LINKABLE
create
  make
feature {NONE}
  make (*a_value*: *INTEGER*)
    -- Create a node with value ‘a_value’.
    do
      value := a_value
    end
feature -- Set links
  set_up (*a_up*: *SKIP_LINKABLE*)
    do
      up := a_up
    end
  set_down (*a_down*: *SKIP_LINKABLE*)
    do
      down := a_down
    end
  set_left (*a_left*: *SKIP_LINKABLE*)
    do
      left := a_left
    end
```

Figure 3: Skip list after promotion of 30-node
set_right (a_right: SKIP_LINKABLE)
do
  right := a_right
end

feature  -- Queries
  value: INTEGER

  up, down, left, right: SKIP_LINKABLE

invariant
  non_circ: left /= Current and right /= Current
end

class
  SKIP_LIST

feature
  minimum: INTEGER
    -- Universal minimal element.
  has (elem: INTEGER): BOOLEAN
    -- Does list contain ‘elem’?
  is_promoted: BOOLEAN
    -- Should a promotion happen?
do
    -- Implementation omitted.
end

search (elem: INTEGER): SKIP_LINKABLE
  -- Rightmost node of the bottom row with value <= ‘elem’.
ensure
  result_exists: Result /= Void
  result_precedes_element: Result.value <= elem
end

insert (elem: INTEGER)
  -- Insert new node with value ‘elem’ into the list.
require
  not has (elem)
local
  ...
ensure
  has (elem)
end

insert_in_row (a_pre, a_node: SKIP_LINKABLE)
  --- Insert node ‘a_node’ after node ‘a_pre’.
require
  nodes_exist: attached a_pre and attached a_node
different_nodes: a_pre /= a_node
local

do

end

\textit{promote} \hspace{1em} (a\_link: \textit{SKIP\_LINKABLE})

\begin{itemize}
\item Possibly promote ‘a\_link’.
\end{itemize}

\textbf{require}
\begin{itemize}
\item node\_exists: \textit{attached} a\_link
\item already\_inserted: \textit{attached} a\_link.left
\end{itemize}

\textbf{local}

\begin{itemize}
\end{itemize}

\textbf{do}

\begin{itemize}
\end{itemize}
end
end