

Pragmatic Type Interoperability

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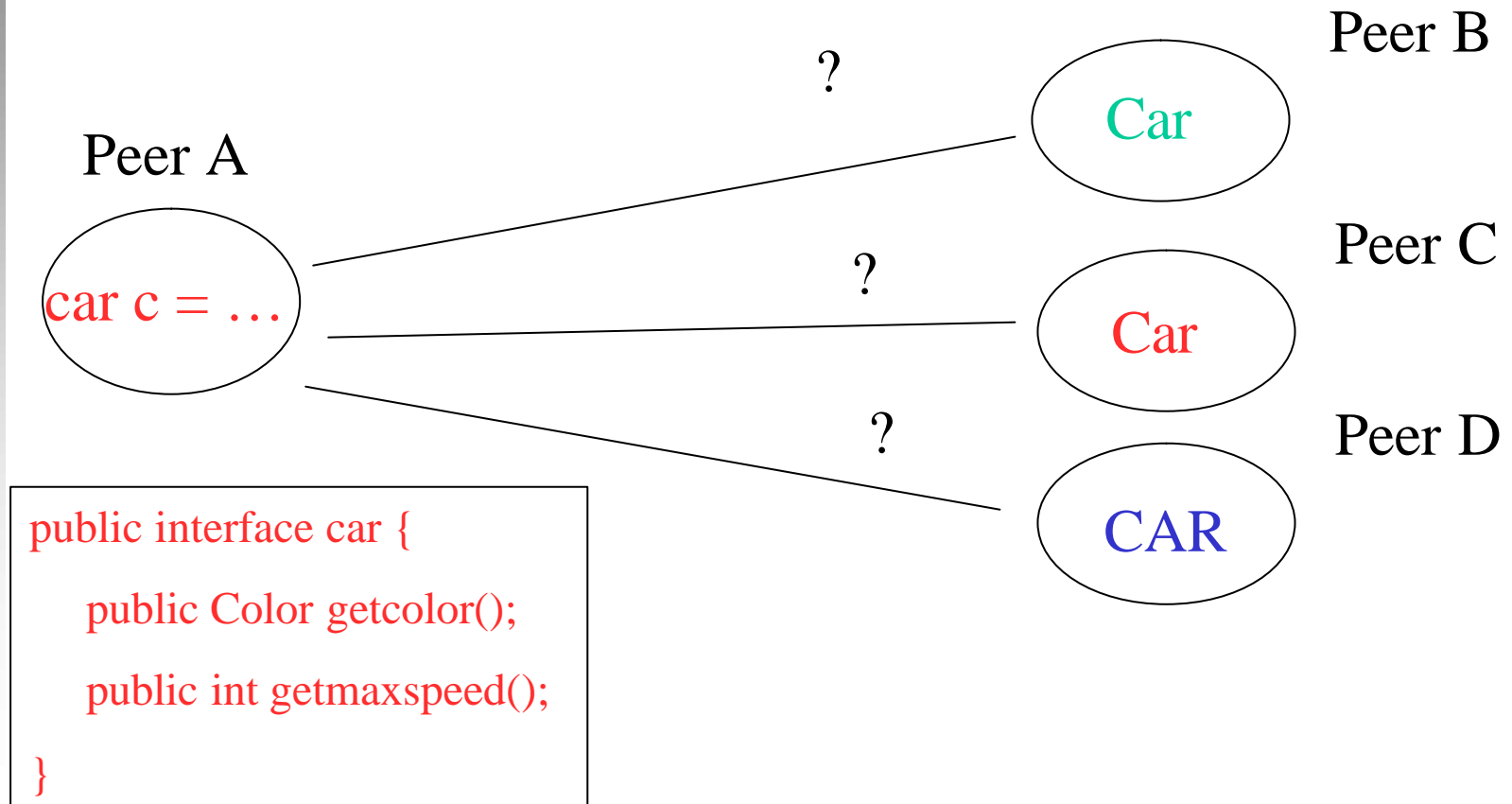
Roadmap

- Context & motivation
- Faces of interoperability
- Object passing semantics & seminal work
- Simplified protocol & implementation issues
- Overhead
- Conclusions & future work

Context & Motivation

- Middleware for peer-to-peer (p2p) object sharing
 - *Borrow/lend* (BL) abstraction, publish/subscribe variant
- Type safety
 - Ideally static within component/processes
- Scalability
 - Abstraction, algorithms
- Interoperability

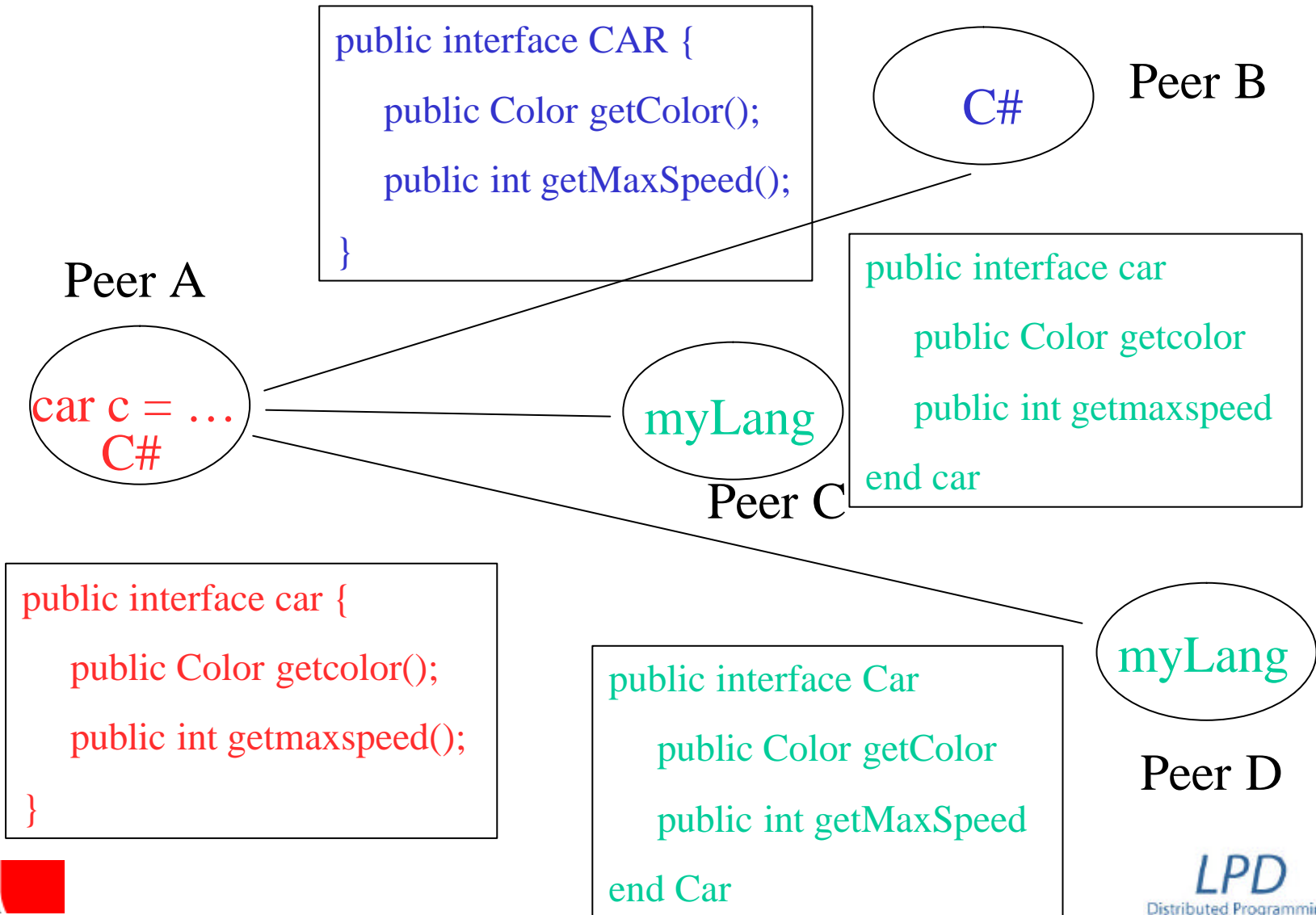
Illustration



Interoperability

- LI: “language interoperability”
 - Same type, different languages
 - Usually explicit agreement
- CI: “component interoperability”
 - Similar type, same language
 - cf. *implicit, structural conformance*
 - Single language centralized settings
- LI & CI
 - Similar type, potentially different languages

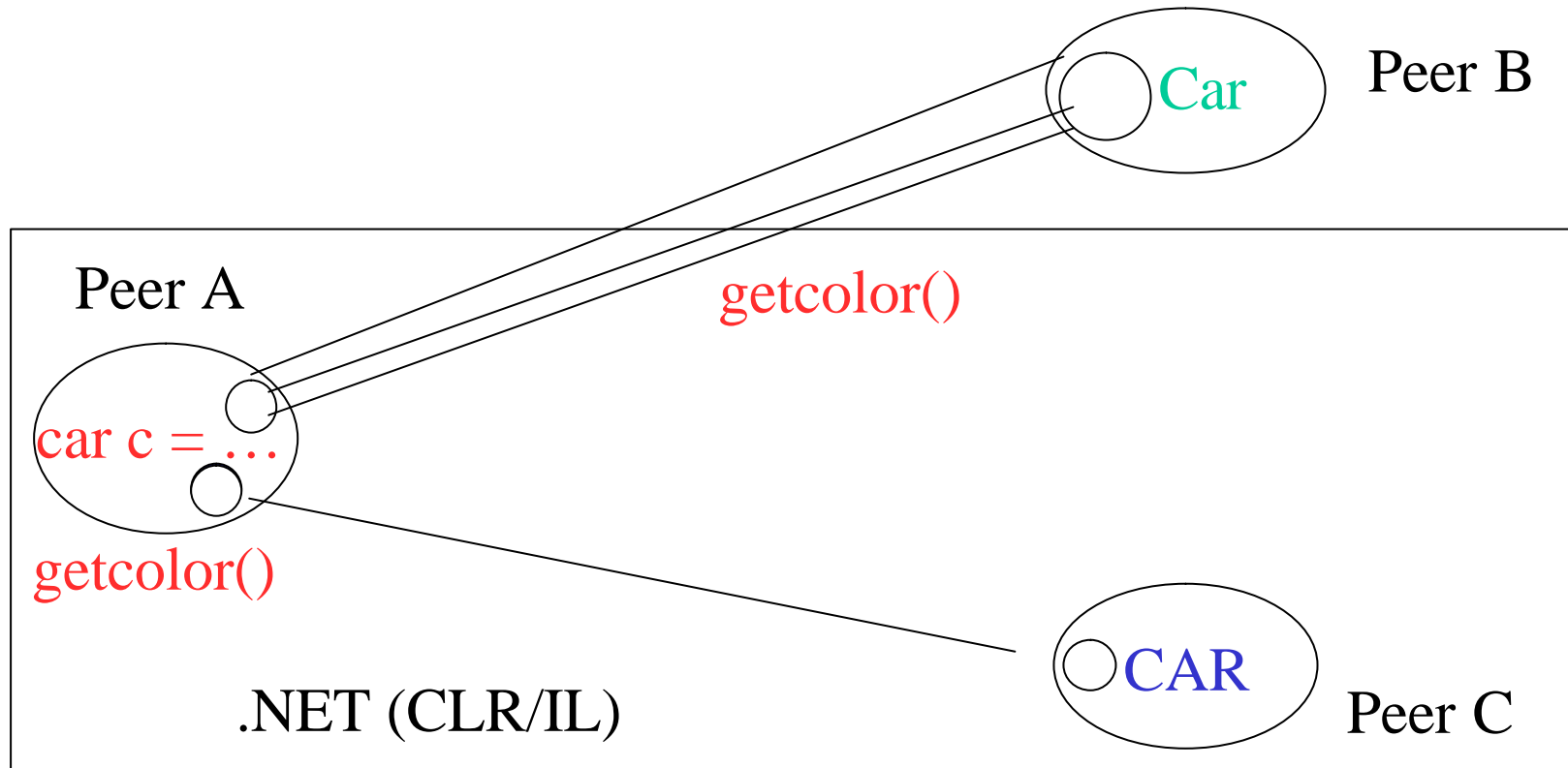
Illustration



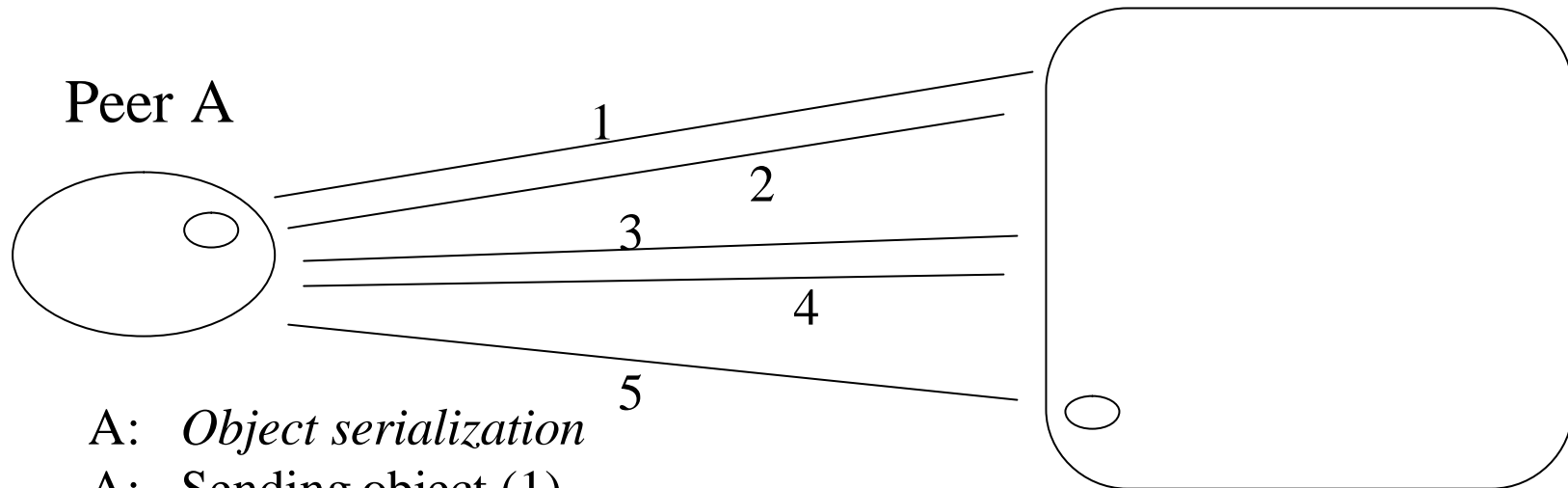
Semantics & Seminal Work

- RS: pass-by-reference semantics
 - Objects invoked remotely (RPC et al.)
 - Proxies and instances of primitive types “transferred”
 - CORBA / .NET (LI), Renaissance [Muckelbauer&Russo] (LI&CI)
- VS: pass-by-value semantics
 - Objects are transferred
 - CORBA (LI limited), .NET (LI)

Illustration



Simple Protocol & Issues (VS)



A: *Object serialization*

A: Sending object (1)

B: Type lookup

B: Asking for new type information (2)

A: Sending *type representation* (3)

B: Applying *type conformance rules*

B: Asking for code, i.e., IL assembly (4)

A: Sending code (5)

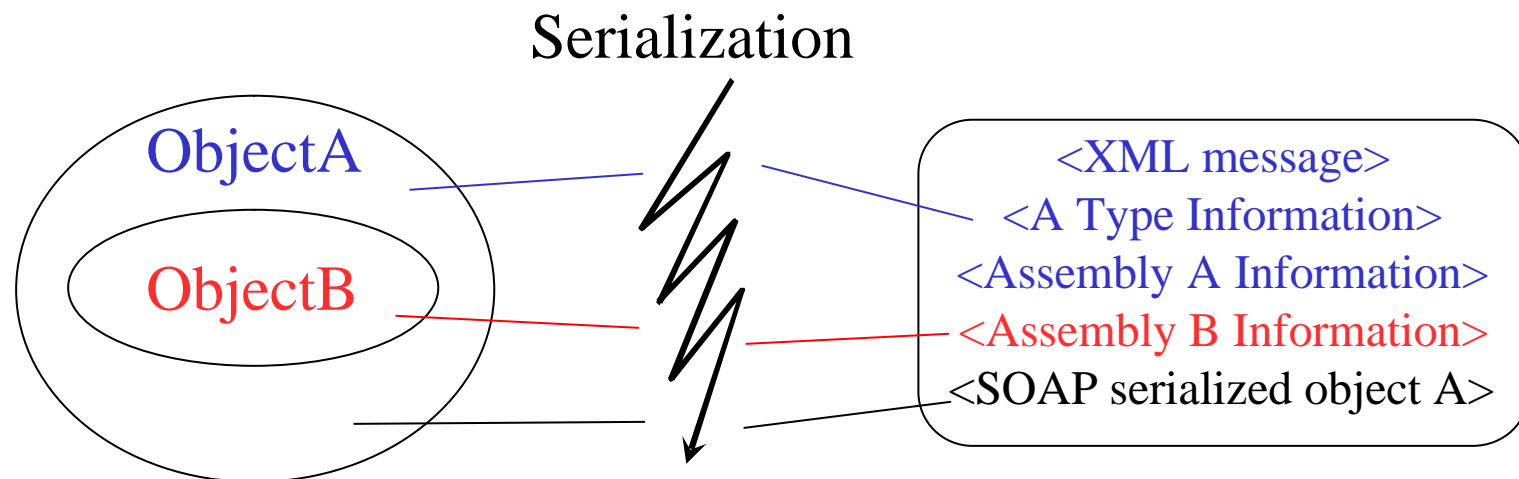
B: Object deserialization

B: *Object access*

Peer B

Object Serialization

- Own XML for *description* of object (type T)
 - Built-in XML serialization only for public fields
 - Binary/SOAP require T at deserialization
- SOAP for object *representation*



Type Representation

- XML serialization with own introspection types
 - Introspection to learn about a type T, but
 - Instances of built-in introspection types (fields, methods, etc.) require type T on target peer

```
public interface CAR {
    public Color getColor();
    public int getMaxSpeed();
}
```

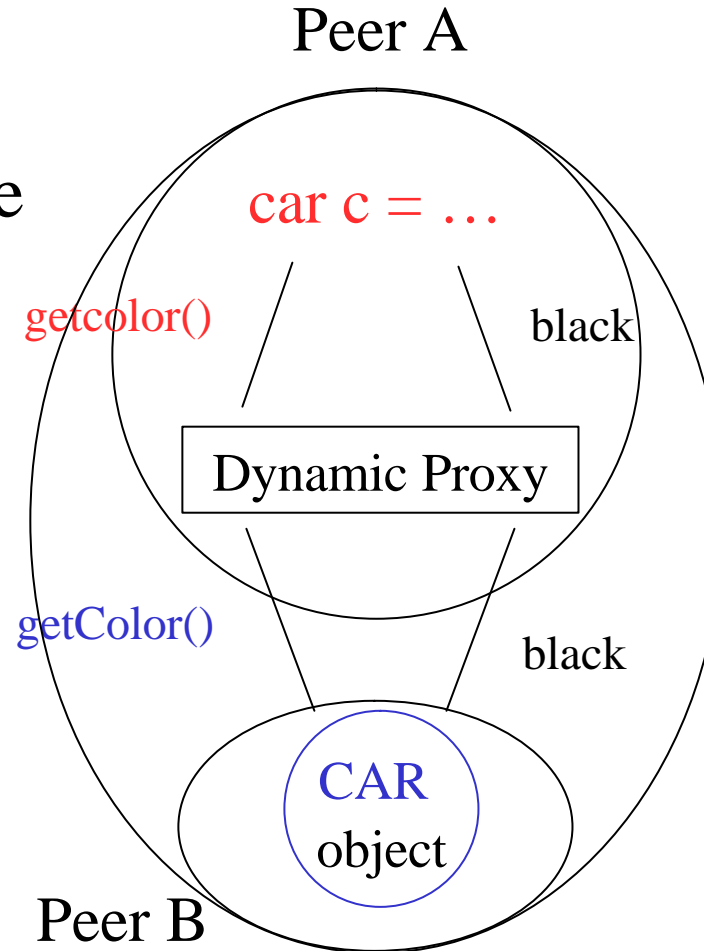
```
<XML message>
<TypeDescription
<classes> <classInfo name="CAR"/></classes>
  <superClasses></superClasses>
  <fields></fields>
  <methods>
    <methodInfo name="getMaxSpeed"
      modifier="public" returnType="int"/>
    ...
  </methods>
</constructors></constructors>
</TypeDescription>
```

Conformance Rules (Informal)

- T conforms to T' explicitly or implicitly
- T implicitly conforms to T' iff
 - Every method in T' (supertypes) finds a method in T (supertypes):
 - Same name (case-insensitive)
 - Parameter types (same order) conform explicitly or implicitly
 - Every field in T' (st) finds a field in T (st):
 - Same name (case-insensitive)
 - Field types conform explicitly or implicitly
- Default (no-argument) constructors

Object Access

- Dynamic proxies
 - Pass-by-reference
 - Pass-by-value



Overhead

Simple types (no recursion)

- Object serialization:
 - Serialization: 16.68 [ms]
 - Deserialization: 1.32 [ms]
- Type representation:
 - Creation and serialization: 6.14 [ms]
 - Deserialization: 2.34 [ms]
- Conformance rules application: 12.66 [ms]
- Object access:
 - Direct invocation: 0.0000142 [ms]
 - Indirect invocation (dynamic proxy): 0.03 [ms]

Conclusions & Future Work

- .NET provides many useful mechanisms
 - Ready-to-use for many scenarios in LI (RS & VS)
 - More effort for achieving CI as not declared goal
 - Reflection costly, need “deeper” integration
- Using conformance (“subtyping”) relationships
 - Object routing
 - Membership management
- Variations on “type conformance”
 - More precise/flexible rules
 - Behavioral conformance for avoiding accidental conformance

Questions and Answers

