Assignment 5: SCOOP principles

ETH Zurich

1 Interpreting a SCOOP program

1.1 Background

The code in listing 1 shows the participants of a crazy office. Note that the BOSS class is the root of this system.

Listing 1: crazy office classes

```
{\bf class} \ BOSS
```

create make

feature

```
evil_supervisor : separate EVIL_SUPERVISOR
nice_supervisor : separate NICE_SUPERVISOR
worker: separate WORKER
```

```
make
```

-- Create supervisors and a worker and use the supervisors to drive the worker. do create evil_supervisor create nice_supervisor create worker print ("boss: I am about to ask the supervisors to do their job.") run (evil_supervisor, nice_supervisor) print ("boss: I am done.")

```
end
```

```
run ( a_evil_supervisor : separate EVIL_SUPERVISOR; a_nice_supervisor: separate
    NICE_SUPERVISOR)
```

-- Use the supervisors to drive the worker.

```
do
```

```
a_evil_supervisor . convince (worker)
a_nice_supervisor . convince (worker)
a_evil_supervisor . convince (worker)
a_nice_supervisor . convince (worker)
```

class EVIL_SUPERVISOR

feature

```
done: BOOLEAN
-- Did I convince a worker?
```

```
convince (a_worker: separate WORKER)
```

-- Convince 'a_worker' that he is not done as soon as he thinks that he is done. require

```
a\_worker.done
```

```
do
```

```
a\_worker.be\_not\_done
done := \mathbf{true}
```

```
print ("evil supervisor: I am done.")
end
```

end

class *NICE_SUPERVISOR*

feature

```
done: BOOLEAN
-- Did I convince a worker?
```

```
convince (a_worker: separate WORKER)
```

— Convince 'a_worker' that he is done as soon as he thinks that he is not done.
require
not a_worker.done
do
a_worker.be_done

```
done := true
```

```
print ("nice supervisor: I am done.")
```

```
end
end
```

```
class WORKER
```

```
create
```

make

feature

```
make
```

-- Create the worker and make him not done.
do
 done := false
 ensure
 not_done: not done
end

```
done: BOOLEAN
-- Do I think that I am done with my task?
```

1.2 Task

Write down one possible output of the program. Does this system terminate (i.e. all processors finish their tasks)?

2 Breakfast Running Time

2.1 Background

Reasoning about the execution times of a concurrent SCOOP program, in the context of break-fast.

2.2 Task

Consider the following SCOOP program being executed on a processor z:

```
bread.cut
toaster.toast
pan.fry
meal.compose
Result := meal.is_cooked and bread.is_delicious
meal.eat
```

The object-processor associations are given as follows: *pan* is handled by processor p, *bread* and *toaster* by processor q, and *meal* by processor r. The call *bread.cut* takes 20 time units until it returns, *toaster.toast* 30 time units, *pan.fry* 20 time units, *meal.compose* 40 time units, *meal.eat* 20 time units. Assume the queries are instantaneous. What is the minimum time for execution of this program? Justify your answer.

3 Baboon Crossing

3.1 Background

This task is adapted from Downey [1] and Tanenbaum [2]. There is a deep canyon somewhere in Kruger National Park, South Africa, and a single rope that spans the canyon. Baboons can cross the canyon by swinging hand-over-hand on the rope, but if two baboons going in opposite directions meet in the middle, they will fight and drop to their deaths. Furthermore, the rope is only strong enough to hold n baboons. If there are more baboons on the rope at the same time, it will break.

3.2 Task

Design and implement a SCOOP synchronization scheme with the following properties:

- Once a baboon has begun to cross, it is guaranteed to get to the other side without running into a baboon going the other way.
- There are never more than n baboons on the rope.
- A continuing stream of baboons crossing in one direction should not bar baboons going the other way indefinitely (no starvation).

References

- [1] Allen B. Downey. The Little Book of Semaphores Second Edition. Green Tea Press, 2005.
- [2] Andrew S. Tanenbaum. Modern Operating Systems (2nd Edition). Prentice Hall, 2001.