



Java and C# in depth

Carlo A. Furia, Marco Piccioni, Bertrand Meyer

C# : concurrency



- C# threads
 - thread implementation
 - sleep and join
 - threads that return values
- Thread synchronization
 - implicit locks and synchronized blocks
 - producer/consumer example
- More efficient concurrency
 - thread pools
 - atomic integers
- Other concurrency models
 - asynchronous programming
 - polyphonic C#



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C# threads

C#'s concurrency model is based on threads

Threads are created by instantiating class **Thread**

- The constructor takes a **ThreadStart delegate** that wraps the method which the thread will execute

Any method can be called with the delegate mechanism

- Unlike Java, any existing class can be used for multi-threaded execution without modifications

In all the examples, assume

```
using System; using System.Threading;
```

A simple class (to be threaded)



```
public class DumbClass {  
  
    private String id;  
  
    public DumbClass(String id) {  
        this.id = id;  
    }  
  
    public void print_id() {  
        // do something  
        Console.WriteLine("This is " + id);  
    }  
}
```

Creating and starting a thread

Create the object with the method the thread will execute

```
DumbClass db = new DumbClass("db");
```

Create a **Thread** object and pass method **print_id** to it using a **ThreadStart** delegate

```
Thread mt = new Thread(  
    new ThreadStart(db.print_id));
```

Start the thread

```
mt.Start();
```

Optionally, wait for it to terminate

```
mt.Join(); // wait until mt terminates  
Console.WriteLine(  
    "The thread has terminated");
```

Putting a thread to sleep



The `Sleep(int t)` **static** method suspends the thread in which it is invoked for `t` milliseconds

```
Thread.Sleep(2000); // suspend for 2 seconds
```

- the timing may be more or less precise according to the real-time guarantees of the executing environment

Threads that return values

Threads can return values using additional delegates

- E.g., to have threads that return strings declare a delegate type:
`public delegate void delForStrings(String s);`
- A class stores a reference to the delegate and activates it when appropriate (to pass values to the caller)

```
public class DullClass {
    private String id;
    // delegate used to return a value when terminating
    private delForStrings d;
    // the constructor must bind the actual method
    public DullClass(String id, delForStrings d)
        { this.id = id;  this.d = d; }
    public void give_id() {
        // call the delegate to return the value id
        if (d != null) { d(id); }
    }
}
```

Creating threads that return values

Define a method to process the information returned by the thread (its signature matches the delegate's)

- for simplicity, we make it static

```
public static void printValueSent(String
s) {
    Console.WriteLine("The thread sent: " + s);
}
```

Create the object with the method the thread will execute and pass the delegate to it

```
DullClass dl = new DullClass("dl",
    new delForString(printValueSent));
```



Creating threads that return values

Create a **Thread** object and pass method **give_id** to it using a **ThreadStart** delegate

```
Thread t = new Thread(  
    new ThreadStart(dl.give_id) );
```

Start the thread

```
t.Start();
```

After executing, it will invoke **printValueSent** through the delegate, which will print the given **id**



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Thread synchronization

Synchronization with locks



The **lock** statement supports synchronization based on locks

- blocks of statements guarded by **lock (o)**
- the lock **o** itself can be any object (including **this**)
- locking/unlocking is implicit when entering/exiting the block
- useful to define critical regions and fine-grained synchronization
- **monitors** are implemented by locking the whole method body on **this**

Synchronization with locks



The **lock** statement supports synchronization based on locks

```
// s must be accessed in mutual exclusion
private int s;

// dict is a read-only object, no concurrency problems
private List<String> dict;

public String decrement_and_lookup() {
    // critical region
    lock(this) { if (s > 0) { s = s - 1; } }
    // non-critical region
    return dict.Item(s);
}
```

Coordination with signals



Locked threads can communicate with **signals**, implemented as static methods of class **Monitor**:

- **Monitor.Wait(o)**: suspend and release the lock on **o** until some thread does a **Pulse(o)** or **PulseAll(o)**
- **Monitor.Pulse(o)**: resume one suspended thread (chosen nondeterministically) waiting on object **o**, which becomes ready for execution when possible
- **Monitor.PulseAll(o)**: resume all suspended threads waiting on object **o**, which become ready for execution when possible
- Analogues of Java's **wait**, **notify**, **notifyAll**

Coordination with events



A more fine-grained (and possibly efficient) coordination uses services of the **WaitHandle** class to coordinate threads

Coordination events are in one of two states: **signaled** and **unsignaled**

- Method **Set** puts an event in the signaled state
 - that is, it issues the signal
- Method **Reset** puts an event in the unsignaled state
 - that is, it cancels the signal



Coordination with events

A more fine-grained (and possibly efficient) coordination uses services of the **WaitHandle** class to coordinate threads

Two main classes implement coordination events

- **AutoResetEvent**
 - automatically resets to unsignaled after being received by one of the waiting threads
- **ManualResetEvent**
 - does not automatically reset, hence it can be received by more than one waiting thread
 - can be reset with method **Reset()**



Coordination with events

Use services of the `WaitHandle` class to coordinate threads

A thread can block **waiting** for an event using some methods of the class

- `WaitHandle.WaitOne()` waits for the event to be signaled (and blocks until then)
- `static WaitHandle.WaitAny(WaitHandle[] e)` waits for any of the events in array `e`.
 - The method returns when an event is received
 - It returns an integer `i`, which is an index within array `e`
 - `e[i]` is the event that has been received
- `static WaitHandle.WaitAll(WaitHandle[] e)` waits for all the events in array `e` to be signaled.

Unlike `Monitor.Wait`, if these wait primitives occur in a **lock** block they do not release the lock while waiting.



The producer-consumer problem

Two threads, the **Producer** and the **Consumer**, work concurrently on a shared **Buffer** of bounded size

The **Producer** puts new messages in the buffer

- if the buffer is full, the Producer must wait until the Consumer takes some messages
- the Producer also signals the last message

The **Consumer** takes messages from the buffer

- if the buffer is empty, the Consumer must wait until the Producer puts some new messages
- the Consumer terminates after the last message

Consistent access to the **Buffer** requires locks and synchronization

One way is to define critical regions when accessing the buffer data structure (with **lock**) and signal events



The main class

```
public class ProducerConsumer {  
  
    public static void main(String[] args) {  
        // create a synchronizer object  
        Synchronizer s = new Synchronizer();  
        // create a buffer of size 3  
        Buffer b = new Buffer(3, s);  
        // create producer and consumer  
        Producer p = new Producer(b, s);  
        Consumer c = new Consumer(b, s);  
        // instantiate threads  
        Thread pT = new Thread(p.produce);  
        Thread cT = new Thread(c.consume);  
        // start them  
        pT.Start();    cT.Start();  
    }  
}
```

Events for synchronization (1/2)



```
using System; using System.Threading;
using System.Collections;
using System.Collections.Generic;

public class Synchronizer {
    private EventWaitHandle takeEvent;
    public EventWaitHandle TakeEvent
        { get { return takeEvent; } }

    private EventWaitHandle giveEvent;
    public EventWaitHandle GiveEvent
        { get { return giveEvent; } }

    private EventWaitHandle endEvent;
    public EventWaitHandle EndEvent
        { get { return endEvent; } }
```

Events for synchronization (2/2)



```
public Synchronizer() {  
    // events initialized to unsignaled state  
    takeEvent = new AutoResetEvent(false);  
    giveEvent = new AutoResetEvent(false);  
    endEvent = new ManualResetEvent(false);  
}  
}
```

- **takeEvent** is an **AutoResetEvent** so it is received by exactly one waiting thread among all those waiting for a take to happen.
- **giveEvent** is an **AutoResetEvent** so it is received by exactly one waiting thread among all those waiting for a give to happen.
- **endEvent** is a **ManualResetEvent** so it is received by all waiting threads: they will all be notified that they can terminate.

The shared Buffer (1/3)



```
public class Buffer {  
  
    public Buffer(int max_size, Synchronizer s) {  
        this.max_size = max_size;  
        this.messages = new Queue<String>();  
        this.s = s;  
    }  
    // buffer of messages, managed as a queue  
    private Queue<String> messages;  
    // maximum number of elements in the buffer  
    private int max_size;  
    // reference to events for synchronization  
    private Synchronizer s;  
}
```

The shared Buffer (2/3)



```
public String take() {
    if (messages.Count == 0) {
        // only one thread receives the event
        WaitHandle.WaitAny(
            // wait until a give occurs
            new WaitHandle[] {s.GiveEvent});
    }
    // now the buffer is not empty
    lock(this) {
        m = messages.Dequeue();
    }
    // signal that a take has occurred
    s.TakeEvent.Set();
    return m;
}
```

The shared Buffer (3/3)



```
public void give(String msg) {
    if (messages.Count == max_size) {
        // only one thread receives the event
        WaitHandle.WaitAny(
            // wait until a take occurs
            new WaitHandle[] {s.TakeEvent});
    }
    // now the buffer has at least an available slot
    lock(this) {
        messages.Enqueue(msg);
    }
    // signal that a give has occurred
    s.GiveEvent.Set();
}
}
```

The Producer (1/2)



```
public class Producer {  
  
    // a reference to the shared buffer  
    private Buffer b;  
  
    // events to synchronize on  
    private Synchronizer s;  
  
    // set the reference to the buffer and synchronizer  
    public Producer(Buffer b, Synchronizer s) {  
        this.b = b;  
        this.s = s;  
    }  
}
```

The Producer (2/2)



```
public void produce() {  
  
    // work for 20 turns  
    for (int i = 0; i < 20; i++) {  
        // put a message in the buffer  
        b.give(i.ToString());  
    }  
    // signal that production has ended  
    s.EndEvent.Set();  
}  
  
}
```

The Consumer (1/2)



```
public class Consumer {  
  
    // a reference to the shared buffer  
    private Buffer b;  
  
    // events to synchronize on  
    private Synchronizer s;  
  
    // set the reference to the buffer and synchronizer  
    public Producer(Buffer b, Synchronizer s) {  
        this.b = b;  
        this.s = s;  
    }  
}
```

The Consumer (2/2)



```
public void consume() {
    // loop as new events arrive, until:
    //      EndEvent is signaled AND b is empty
    while ( WaitHandle.WaitAny( new WaitHandle[]
        {s.GiveEvent, s.EndEvent}) != 1
        || !b.Empty ) {
        string m = b.take();
        Console.WriteLine(
            "Consumer has consumed message " + m );
    }
}
}
```



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More efficient concurrency

Concurrency and performance



Thread creation is time-consuming

- massive thread creation can annihilate responsiveness
- C#'s solution: thread pools

Lower-level primitives are available

- **Mutex** class for mutexes
 - **less** efficient than monitors and **lock** (unlike Java)
- **Interlocked static** class
 - atomic operations on integers

Tip: don't forget the efficiency/abstraction trade-off



Thread pools

Thread pools are an efficient way of running multi-threaded applications

- maintain a pool of worker threads
- when a client requests a new task to run, preempt one of the available worker threads and assign it to the task
- no creation overhead upon task invocation

C#'s static class `System.Threading.ThreadPool`

- `QueueUserWorkItem(WaitCallback w, Object o)`: schedule delegate `w` for execution by a worker thread, when possible; `o` is passed as argument to `w`.

Thread pool thread creation

Create a wrapper delegate for each method to be threaded

- In the Producer/Consumer example:

```
public static void Main(string[] args) {
    Producer p = new Producer(b, s);
    Consumer c = new Consumer(b, s);
    ThreadPool.QueueUserWorkItem(new
        WaitCallback(consuming), c);
    ThreadPool.QueueUserWorkItem(new
        WaitCallback(producing), p);
}

public static void consuming(object o)
{ ((Consumer) o).consume(); }

public static void producing(object o)
{ ((Producer) o).produce(); }
```

There's an undesirable side-effect with this code as is.
What is it?

Thread pool thread creation

Create a wrapper delegate for each method to be threaded

- In the Producer/Consumer example:

```
public static void Main(string[] args) {
    Producer p = new Producer(b, s);
    Consumer c = new Consumer(b, s);
    ThreadPool.QueueUserWorkItem(new
        WaitCallback(consuming), c);
    ThreadPool.QueueUserWorkItem(new
        WaitCallback(producing), p);
}
```

There's an undesirable side-effect with this code as is.
What is it?

- **Main** terminates after invoking **QueueUserWorkItem**; hence the **ThreadPool** object is deallocated and the worker threads forcefully terminated!

Interlocked class



C#'s implementation of atomic operations on integers

```
// shared variable
```

```
int s;
```

```
...
```

```
// this is equivalent to an atomic s++
```

```
Interlocked.Increment(ref s);
```

```
...
```

```
// this is equivalent to an atomic s--
```

```
Interlocked.Decrement(ref s);
```



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Other concurrency models:
Asynchronous programming

Concurrency and correctness



Programming thread-safe data structures is error-prone

- Thread-safe collections are available since C# 4.0
- Current collections provide a **SyncRoot** object for synchronization

Threads and monitors are too general for straightforward parallel computation

- C#'s solution: asynchronous methods

Tip: don't forget the efficiency/abstraction trade-off

Asynchronous programming



C# 5.0 introduced simple mechanisms to have methods execute asynchronously and wait for one another.

The model is based on asynchronous methods:

```
async Task<T> DoAsync ()
```

- **DoAsync** may execute asynchronously from its clients
- In turn, its clients can wait for **DoAsync**'s to complete (and only then access its result).

(The class **Task** can also be used independent of asynchronous methods, mostly to introduce forms of data-bound parallelism.)

Asynchronous methods



`async Task<T> DoAsync ()`

Asynchronous methods:

- Are declared as such with the keyword `async`
- Can have only specific return types:
 - `Task<T>` for methods returning values of type `T`
 - `Task` for methods returning no values
 - `void` for methods returning no values used as event handlers
- Cannot have `ref` or `out` arguments (there's no way to "wait" for those)
- By convention, have name ending in "`Async`"
- Can wait for other asynchronous methods to complete using the `await` instruction in their bodies.

```
async Task<T> DoAsync ()
```

When an asynchronous method `DoAsync` executes an `await`:

- Control **may** return to the caller (the compiler/runtime decides if a context switch is worth the cost)
- The caller will be able to retrieve the result later when available, after `awaiting`
- No new thread is created: the asynchronous computation uses the thread executing `DoAsync`

The result obtained when `awaiting` for an asynchronous method with return type `Task<T>` has type `T`.

Asynchronous programming: example



Write a method `AvgAgesAsync` that computes the average age of the population of several cities.

The data for each city is accessible remotely using a library method:

```
async Task<List<int>> GetAgesAsync (String city)
```

A call return a list of ages, one for each person of the city.

Calls to `AvgAgesAsync` may take time, but can be executed asynchronously:

1. First, the client start the asynchronous computation:

```
Task<double> t = AvgAgesAsync (listOfCities) ;
```

2. Now, the client can do other stuff while `AvgAgesAsync` executes in parallel.
3. Eventually, the client will get the final results with a call:

```
double avg = await t;
```

Asynchronous programming: example



```
async Task<double> AvgAgesAsync (List<String> cities)
{
    int i = 0, pop = 0; double avg = 0;
    foreach (String c in cities) {
        // wait for results from GetAgesAsync
        // (but AvgAgesAsync's caller needn't block)
        List<int> v = await GetAgesAsync(c);
        avg = // new average, from old one
            ((avg*pop) + v.Sum()) / (pop + v.Count);
        pop += v.Count; // new total population
        i++; // one more city done
        Console.WriteLine(
            "Done {0}% of cities. Current average: {1}",
                (i/cities.Count*100), avg);
    } return avg;
}
```



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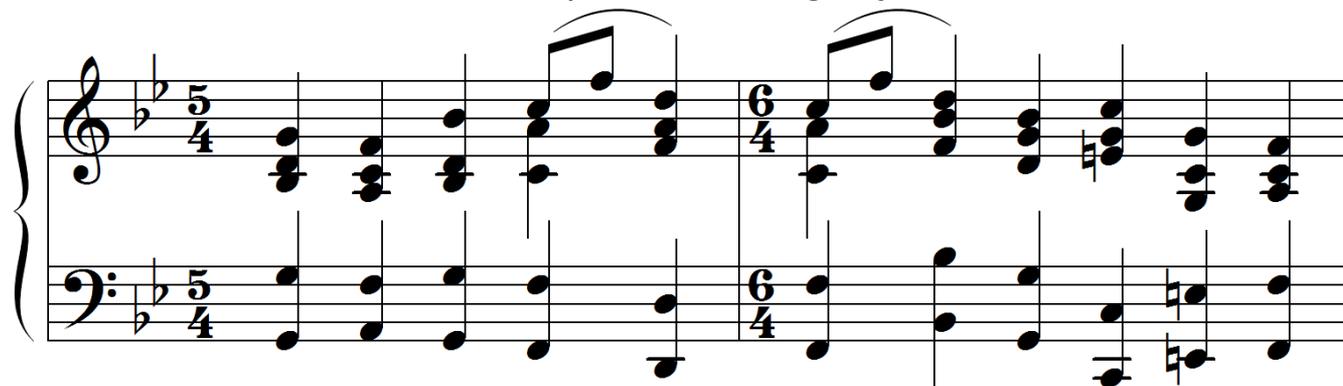
Other concurrency models:
Polyphonic C#

Introducing Polyphonic C#



- **Polyphonic C#** is an extension of C# with a few high-level primitives for concurrency
 - not part of .NET framework
 - based on join calculus (Fournet & Gonthier, 1996)
 - taken up by Microsoft's **C ω** project
 - **JoinJava** is a similar extension for Java
- Based on two basic notions
 - Asynchronous methods
 - Chords

(M. Mussorgsky, Pictures at an exhibition)



Polyphonic asynchronous methods



Calls to asynchronous methods return immediately without returning any result

- The callee is scheduled for execution in a different thread
- similar to sending a message or raising an event
- declared using **async** keyword instead of **void**

```
public async startComputation () {  
    // computation  
}
```

- asynchronous methods do not return any value and cannot have **ref** or **out** arguments

Chords: syntax



A chord is an extension of the notion of method definition

- The signature of a chord is a collection of (traditional) method declarations joined by `&`
- The body of a chord is all similar to the body of a traditional method

```
public int get() & public async put(int i) {  
    return i;  
}
```

- Within a chord:
 - at most one method can be non-asynchronous
- Within a class:
 - the same method can appear in more than one chord
- (We do not discuss additional rules for inheritance and overloading)



A chord is only executed once all the methods in its signature have been called

- Calls are buffered until there is a matching chord
 - the implicit buffer supports complex synchronization patterns with little code (see Producer/Consumer later)
- If multiple matches are possible, nondeterminism applies
- Execution returns a value to the only non-asynchronous method in the chord (if any)

Chords semantics: example



```
public class Buffer() {
    public int get() & public async put(int i)
        { return i; }
}

...
Buffer b = new Buffer();
b.put("okey")
Console.WriteLine(b.get()); // prints "okey"
b.put("okey"); b.put("dokey");
    // prints "okeydokey" or "dokeyokey"
Console.WriteLine(b.get() + b.get());
b.get(); // blocks until some other thread calls put
```

Producer/Consumer with chords (1/3)



```
public class Buffer {  
  
    public void give(String s) & async available(int a) {  
        if (a == 1) {  
            // just one slot available and giving: become full  
                full();  
        } else {  
            // more than one slot available and giving:  
            // enable more giving  
                available(a - 1);  
        }  
        // buffer message for takes  
        inBuffer(s);  
    }  
}
```

Producer/Consumer with chords (2/3)



```
public String take() & async inBuffer(String s) &
    async full() {
    // full and taking: one slot becomes available
    available(1);
    // return message in queue
    return s;
}
```

```
public String take() & async inBuffer(String s) &
    async available(int a) {
    // not full: one more slot becomes available
    available(a + 1);
    // return message in queue
    return s;
}
```

Producer/Consumer with chords (3/3)



```
// constructor
public Buffer(int capacity) {
    available(capacity);
}
}
```

Note: unlike in the examples we developed with locks, here there is no guarantee of ordered retrieval because any message in the implicit buffer can be retrieved at any time