Assignment 4: SCOOP principles

ETH Zurich

1 Interpreting a SCOOP program

1.1 Background

The code in listing 1 shows the participants of a crazy office. Note that the BOSS class is the root of this system.

Listing 1: crazy office classes

```
class BOSS
create
 make
feature
  evil\_supervisor: separate EVIL\_SUPERVISOR
  nice_supervisor: separate NICE_SUPERVISOR
  worker: separate WORKER
  make
      — Create supervisors and a worker and use the supervisors to drive the worker.
   do
     create evil_supervisor
     create nice_supervisor
     create worker
     print ("boss: I am about to ask the supervisors to do their job.")
     run ( evil_supervisor , nice_supervisor )
     print ("boss: I am done.")
   end
  run (a_evil_supervisor: separate EVIL_SUPERVISOR; a_nice_supervisor: separate
      NICE_SUPERVISOR)
      — Use the supervisors to drive the worker.
   do
      a\_evil\_supervisor . convince (worker)
      a_nice_supervisor.convince (worker)
      a_evil\_supervisor . convince ( worker)
      a_nice_supervisor.convince (worker)
     if (a_evil_supervisor.done and a_nice_supervisor.done) then
       print ("boss: The supervisors are done.")
     end
   end
end
```

```
{\bf class}\ EVIL\_SUPERVISOR
feature
 done: BOOLEAN
    -- Did I convince a worker?
  convince (a_worker: separate WORKER)
      -- Convince 'a_worker' that he is not done as soon as he thinks that he is done.
   require
      a\_worker.done
   do
      a\_worker.be\_not\_done
      done := \mathbf{true}
      print ("evil supervisor: I am done.")
end
class NICE_SUPERVISOR
feature
  done: BOOLEAN
   -- Did I convince a worker?
  convince (a_worker: separate WORKER)
      -- Convince 'a_worker' that he is done as soon as he thinks that he is not done.
   require
     not a_worker.done
   do
      a\_worker.be\_done
      done := \mathbf{true}
      print ("nice supervisor: I am done.")
   end
end
class WORKER
create
 make
feature
 make
      -- Create the worker and make him not done.
   do
      done := false
   ensure
      not_done: not done
   end
  done: BOOLEAN
   -- Do I think that I am done with my task?
```

1.2 Task

Write down one possible output of the program. Does this system terminate (i.e. all processors finish their tasks)?

1.3 Solution

The system terminates. One of the possible outputs is:

- 1. boss: I am about to ask the supervisors to do their job.
- 2. nice supervisor: I am done.
- 3. worker: I am done.
- 4. evil supervisor: I am done.
- 5. worker: I am not done.
- 6. nice supervisor: I am done.
- 7. worker: I am done.
- 8. evil supervisor: I am done.
- 9. worker: I am not done.
- 10. boss: The supervisors are done.
- 11. boss: I am done.

Variations of the above output are given by the fact that a worker can print its message before the supervisor and the other way around. The remaining orderings are predefined by the program.

2 Breakfast Running Time

2.1 Background

Reasoning about the execution times of a concurrent SCOOP program, in the context of breakfast.

2.2 Task

Consider the following SCOOP program being executed on a processor z:

```
bread.cut
toaster.toast
pan.fry
meal.compose
Result := meal.is_cooked and bread.is_delicious
meal.eat
```

The object-processor associations are given as follows: pan is handled by processor p, bread and toaster by processor q, and meal by processor r. The call bread.cut takes 20 time units until it returns, toaster.toast 30 time units, pan.fry 20 time units, meal.compose 40 time units, meal.eat 20 time units. Assume the queries are instantaneous. What is the minimum time for execution of this program? Justify your answer.

2.3 Solution

The bread and toaster must run in sequence, taking 50 time units. The pan and the first meal take 20 and 40 time units respectively. All 3 of these times are run in parallel, so their combined running time is the maximum, or 50 time units. The program then synchronizes at the **Result** line, waiting for the response of the meal and bread. There is an additional 20 time unit delay at the end.

The total running time is then 70 time units.

3 Baboon Crossing

3.1 Background

This task is adapted from Downey [1] and Tanenbaum [2]. There is a deep canyon somewhere in Kruger National Park, South Africa, and a single rope that spans the canyon. Baboons can cross the canyon by swinging hand-over-hand on the rope, but if two baboons going in opposite directions meet in the middle, they will fight and drop to their deaths. Furthermore, the rope is only strong enough to hold n baboons. If there are more baboons on the rope at the same time, it will break.

3.2 Task

Design and implement a SCOOP synchronization scheme with the following properties:

- Once a baboon has begun to cross, it is guaranteed to get to the other side without running into a baboon going the other way.
- There are never more than n baboons on the rope.
- A continuing stream of baboons crossing in one direction should not bar baboons going the other way indefinitely (no starvation).

3.3 Solution

A solution can be found in the SCOOP example directory, which is part of the EiffelStudio installation.

References

- [1] Allen B. Downey. The Little Book of Semaphores Second Edition. Green Tea Press, 2005.
- [2] Andrew S. Tanenbaum. Modern Operating Systems (2nd Edition). Prentice Hall, 2001.