



Java and C# in depth

Carlo A. Furia, Marco Piccioni, Bertrand Meyer

Java: framework overview
and in-the-small features



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Java: framework overview



What's in a name

Initially was “Oak” (James Gosling, 1991), then “Green”

- Ruled out by the trademark lawyers

Twelve people locked in a room together with a “naming consultant”

- “How does this thing make you feel?”
- “What else makes you feel that way?”

After listing and sorting, 12 names were sent to the lawyers

- #1 was “Silk”
- Gosling’s favorite was “Lyric” (#3)
- “Java” was # 4

Version 1.0: 1995, latest stable version: 7 Update 51 (14.1.14)

Coming next (Java SE 8, 18.3.2014):

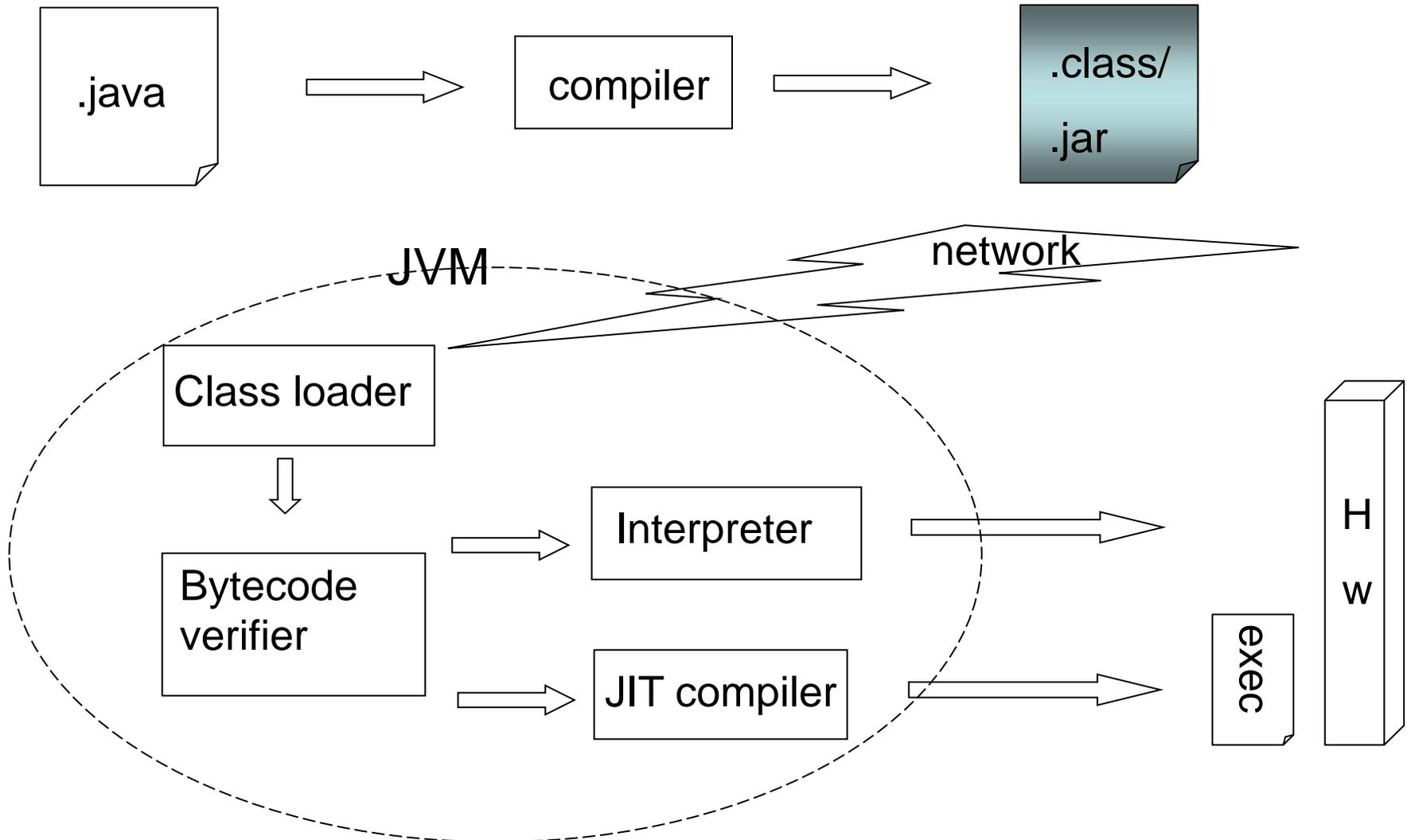
- lambda expressions (closures)
- embedded JavaScript

Java platform goals



- Write Once, Run Anywhere
- Built-in security
- Automatic memory management
- API + documentation generation
- Object-Oriented
- Familiar C/C++ syntax

Write once, run anywhere





- Intermediate format resulting from Java compilation
- Instruction set of an architecture that
 - is stack-oriented (no registers)
 - provides capability (object access rights)
- 1 bytecode instruction = 1 byte
- Executed by any platform-specific Virtual Machine (VM)

Bytecode format



- JVM loads class file → gets a stream of bytecodes
- One bytecode instruction: **opcode** + ≥ 0 **operands**
- Each opcode is associated with a **mnemonic**
 - 03 → `iconst_0` // pushes int 0 on stack
 - 3b → `istore_0` // pops int from stack to local in pos 0
 - 84 00 01 → `iinc 0, 1` // increments local in pos 0 by 1
 - 1a → `iload_0` // pushes int from local in pos 0 on stack
 - 05 → `iconst_2` // pushes int 2 on stack
 - 68 → `imul` // pops 2 int values, multiplies them and puts the result on the stack

Example of bytecode translation



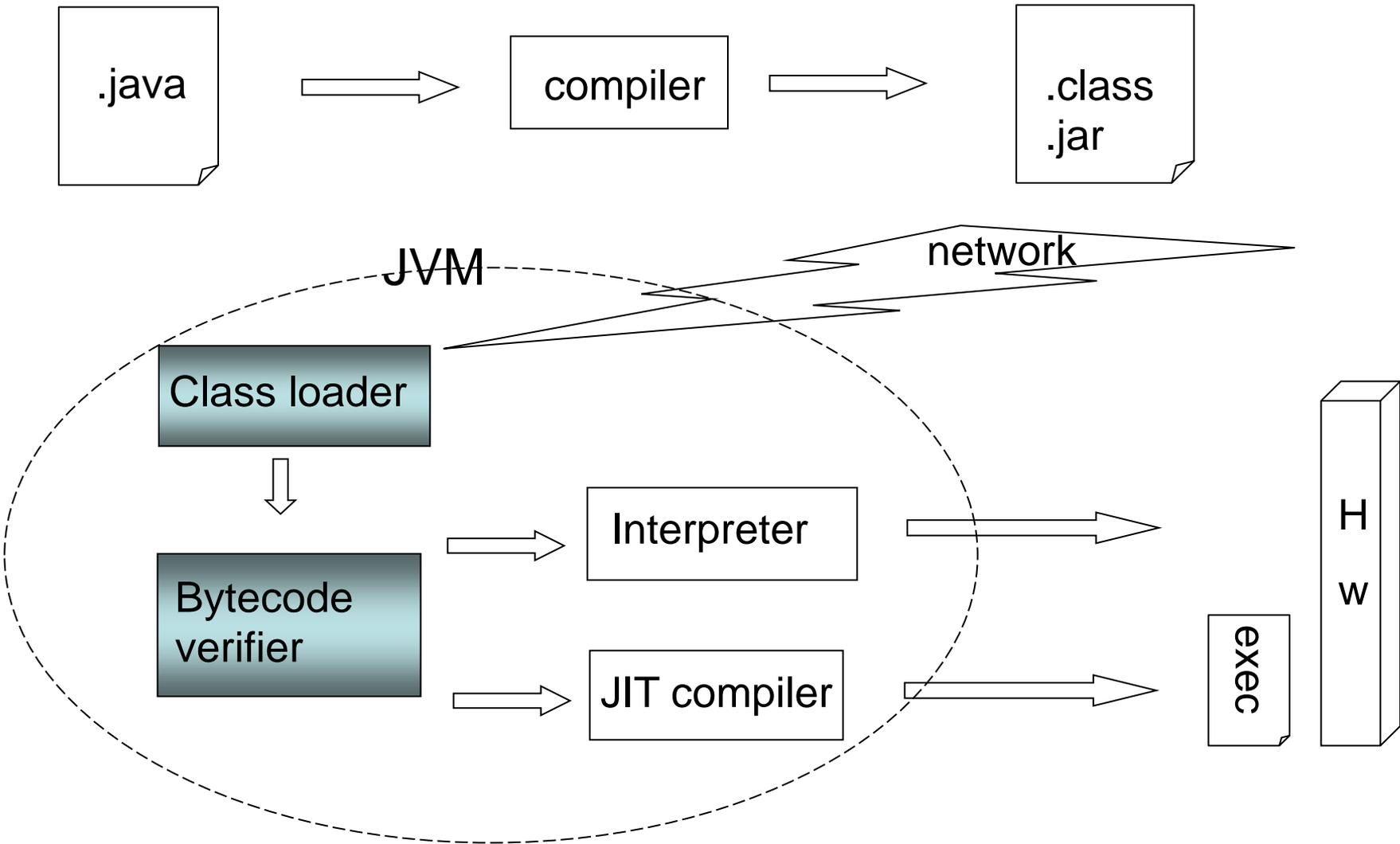
```
class SimpleMath{
    byte inflexible_add() {
        byte x = 2;
        byte y = 2;
        byte z = (byte) (x + y);
        return z;
    }
}
```

Bytecode example



Opcode mnemonics	Meaning
iconst_2	push an integer constant 2 into the stack
istore_1	pop into local in pos 1 (x)
iconst_2	push an integer constant 2 into the stack
istore_2	pop into local in pos 2 (y)
iload_1	push x into the stack
iload_2	push y into the stack
iadd	sum the two top values on the stack and push the result
int2byte	convert result into byte
istore_3	pop into local in pos 3 (z)
iload_3	push z into the stack
ireturn	return result (z)

JVM overview





- No pointers, no explicit memory de-allocation
- Checked type casts (at compile time and runtime)
- Enforced array bounds (at runtime)
- Security APIs
 - SecurityManager (standard security)
 - XML digital signature, Public Key Infrastructure, cryptographic services, authentication

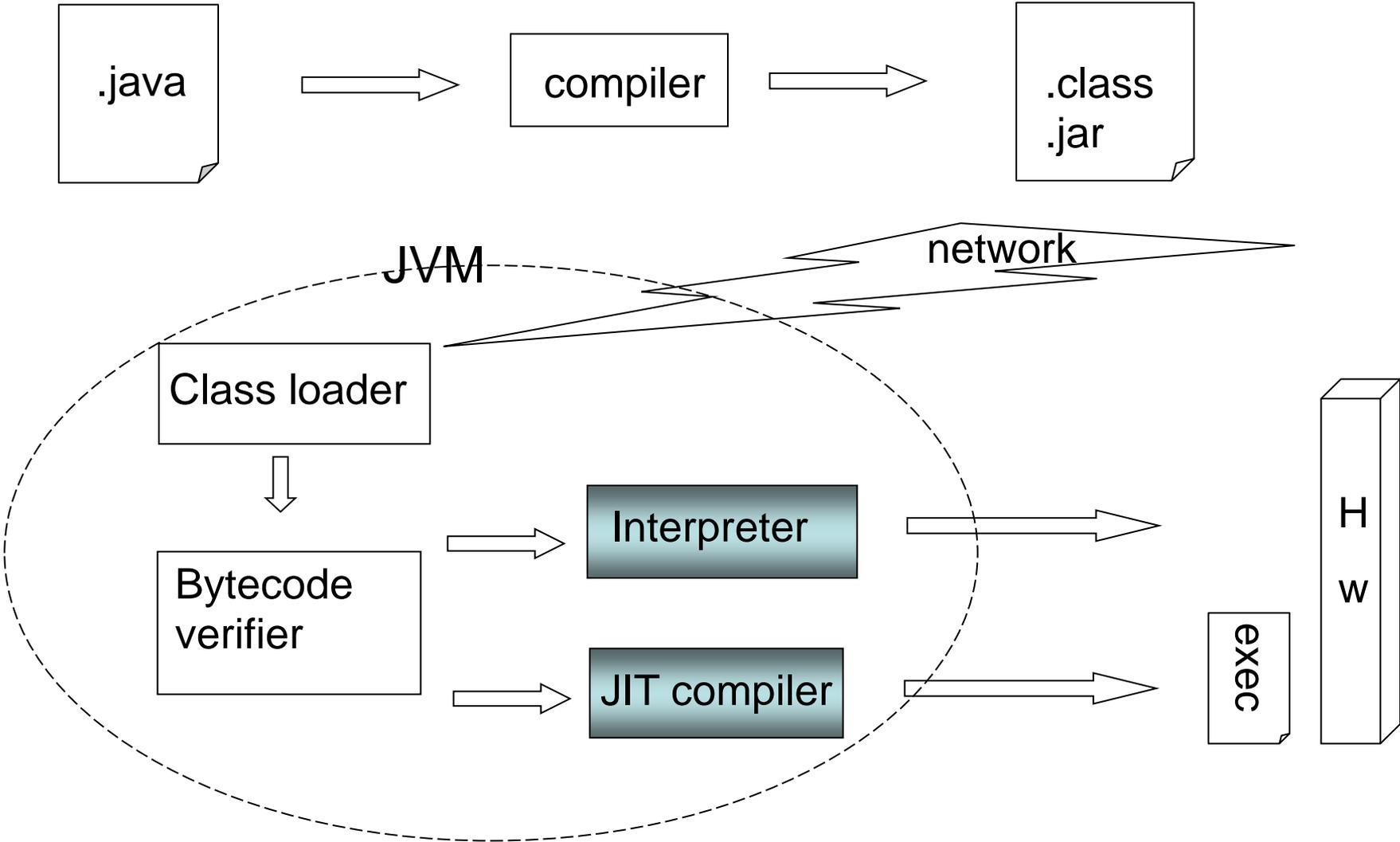


- Take care of files and file systems
- Locate libraries and dynamically load classes
- Partition classes into realms (e.g. local machine, local network, all the rest) and restrict what they can do



- Verifier checks bytecode using a “theorem prover”
 - Branches always to valid locations
 - Data always initialized
 - Types of parameters of bytecode instructions always correct
 - Data and methods access checked for visibility
 - Arbitrary bit patterns cannot get used as an address
 - No operand stack overflows and underflows

JVM: code generation



Code generation: HotSpot



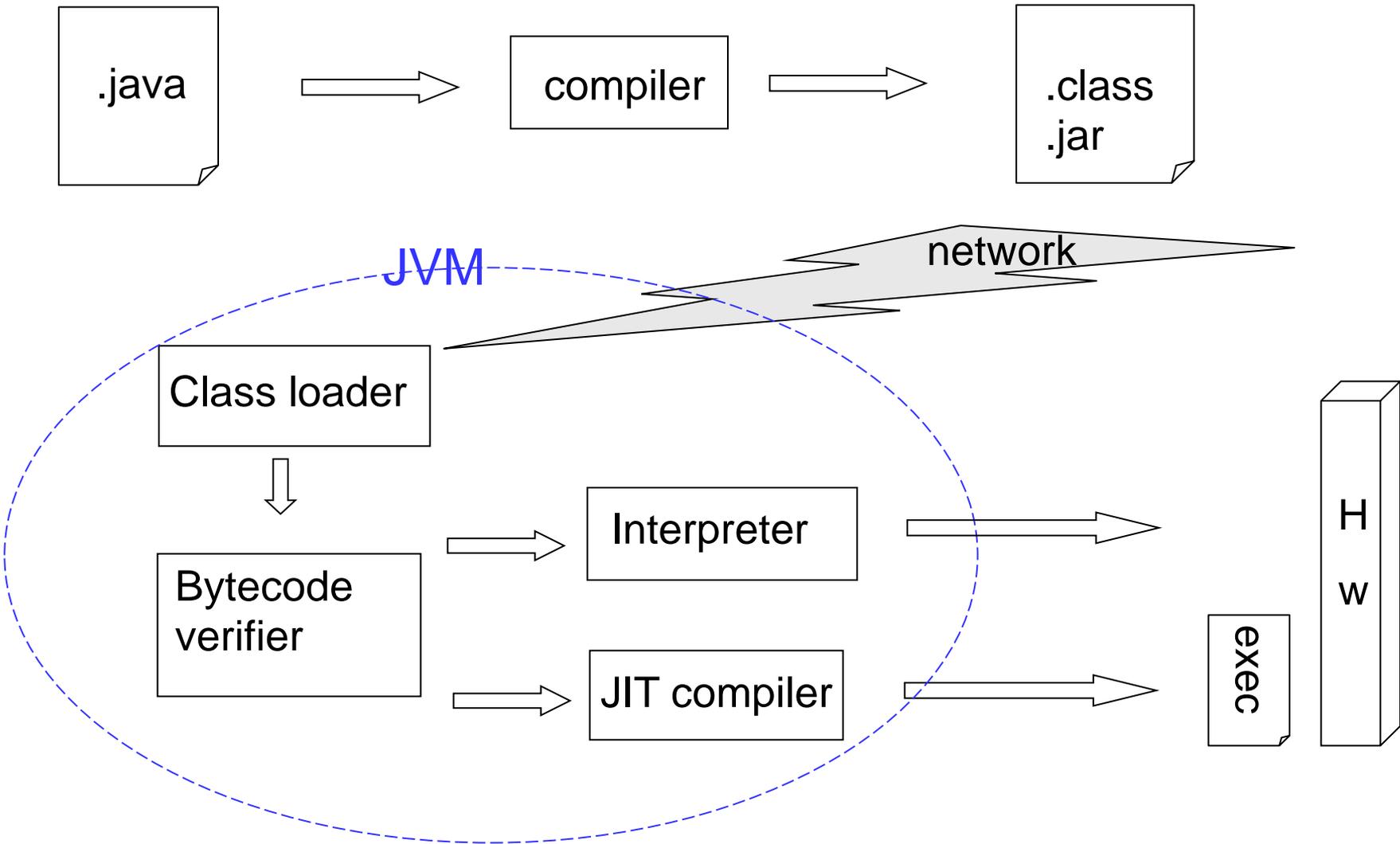
- The interpreter is the software CPU of the JVM
 - Examines each bytecode and executes a unique native procedure
 - No native code is produced
- A JIT “compiler” converts the bytecode into native code just before running it
 - Keeps a log (cache) of the native code that it has to run to execute each bytecode
 - May optimize substituting often occurring short sets of instructions (“hot spots”) with shorter/faster ones
 - Like the back-end of a traditional compiler, the java compiler being the front-end
- HotSpot is the default SUN JVM since 2000



HotSpot client and server

- HotSpot client VM
 - For platforms typically used for client applications (e.g. GUI)
 - Tuned for reducing start-up time and memory footprint
 - Invoked by using *-client* when launching an app
- HotSpot server VM
 - For all platforms
 - Tuned for max program execution speed
 - Invoked by using *-server* when launching an app
- Both use an interpreter to launch applications, and an adaptive compiler optimizing code hot spots
- They use different code inline policies and heap defaults

JVM Overview





- Automated exception handling
 - Provides “root cause” debugging info for every exception
- Responsible for garbage collection
- Ships as JRE (VM + libraries)
- Can have other languages run on top of it, e.g.
 - JRuby (Ruby)
 - Rhino (JavaScript)
 - Jython (Python)
 - Scala
- From 6.0 scripting languages can be mixed with Java code



- Compile

```
javac MainClass.java
```

- Execute

```
java MainClass
```

- Generate documentation

```
javadoc MainClass.java
```

- Generate an archive from `.class` files in current dir

```
jar cf myarchive.jar *.class
```



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Java: in-the-small language features

Encoding and formatting



- Uses unicode as encoding system: www.unicode.org
- Free format
 - Blanks, tabs, new lines, form feeds are only used to keep tokens separate
- Comments
 - Single line: `//Single line comment`
 - Multiple lines: `/* non-nested, multi-line
comment*/`
 - Javadoc comment: `/** processed by javadoc */`



- No restriction on length
- Case sensitive
- Cannot start with a digit
- Cannot include / or -
- Cannot be a keyword



Meta-data about programs

- Compiler flags
e.g: `@Deprecated`, `@Override`, `@SuppressWarnings`
- Information that can be used for compilation (or other forms of code analysis)
e.g.: `@Inherited`, application-defined such as `@RevisionId`
- Some runtime processing
e.g.: application-defined

Keywords



<code>abstract</code>	<code>double</code>	<code>int</code>	<code>super</code>
<code>boolean</code>	<code>else</code>	<code>interface</code>	<code>switch</code>
<code>break</code>	<code>extends</code>	<code>long</code>	<code>synchronized</code>
<code>byte</code>	<code>final</code>	<code>native</code>	<code>this</code>
<code>case</code>	<code>finally</code>	<code>new</code>	<code>throw</code>
<code>catch</code>	<code>float</code>	<code>package</code>	<code>throws</code>
<code>char</code>	<code>for</code>	<code>private</code>	<code>transient</code>
<code>class</code>	<code>(goto)</code>	<code>protected</code>	<code>try</code>
<code>(const)</code>	<code>if</code>	<code>public</code>	<code>void</code>
<code>continue</code>	<code>implements</code>	<code>return</code>	<code>volatile</code>
<code>default</code>	<code>import</code>	<code>short</code>	<code>while</code>
<code>do</code>	<code>instanceof</code>		

- Literals `null`, `true`, `false` are also reserved

Operators



- Access, method call: `.`, `[]`, `()`
- Postfix: `expr++`, `expr--` (R to L)
- Other unary: `++expr`, `--expr`, `+`, `-`, `~`, `!`, `new`, `(aType)`
- Arithmetic: `*`, `/`, `%`
- Additive: `+`, `-`
- Shift: `<<`, `>>`, `>>>`
- Relational: `<`, `>`, `<=`, `>=`, `instanceof`
- Equality: `==`, `!=`
- Logical (L to R): `&`, `^`, `|`, `&&`, `||`
- Ternary: `condition ? (expr1) : (expr2)` (R to L)
- Assignment: `=`, `+=`, `-=`, `*=`, `/=`, `%=`, `&=`, `^=`, `|=`, `<<=`, `>>=`, `>>>=`
- **Precedence**: from top to bottom
- **Tip**: don't rely too much on precedence rules: use parentheses

Type system basics



- Primitive types
 - `boolean, byte, short, int, long, char, float, double`
- Reference types
 - `class, interface, []`
- `null`
- Automatic widening conversions (no precision loss)
 - `byte` to `short` to `int` to `long`
 - `char` to `int`, `int` to `double`, `float` to `double`
- Automatic widening conversions (possible precision loss)
 - `int` to `float`, `long` to `float`, `long` to `double`
- A cast is required for narrowing conversions
 - `int i = 3; long j = 5; i = (int) j`

Widening conversions with precision loss

```
float g(int x) {  
    return x;  
}  
  
...  
int i = 1234567890;  
float f = g(i);  
System.out.println(i - (int) f)  
// output: -46  
...
```

Wrapper types and autoboxing



- For each primitive type there is a wrapper type
 - **Boolean, Byte, Short, Integer, Long, Character, Float, Double**
- Starting from 5.0, autoboxing provides automatic conversions between primitive and wrapper types
- Pro: reduces code complexity
- Cons: not efficient, sometimes unexpected behavior

Some surprises of autoboxing



```
new Integer(7).equals(7) // true
```

```
new Long(7).equals(7) //false. True if equals(7L)
```

```
new Integer(7).equals(new Long(7)) // false
```

```
new Integer(7) == 7 // true
```

```
new Long(7) == 7 // true
```

```
new Integer(7) == new Long(7) // compiler error
```

Control flow: conditional branch



Same syntax as in C/C++

```
if (booleanExpr)
{
    // do something
}
else // else is optional
{
    // do something else
}
```

Control flow: loops



```
while (booleanExpr)
{
    // execute body
    // until booleanExpr becomes false
}
```

```
do
{
    // execute body (at least once)
    // until booleanExpr becomes false
}
```

```
while (booleanExpr) ;
```

Control flow: **for** loop



```
for (int i=0; i < n; i++)  
{  
    // execute loop body n times  
}
```

// equivalent to the following

```
int i=0;  
while (i < n)  
{  
    // executes loop body n times  
    i++;  
}
```

Control flow: enhanced **for** loop



Introduced in Java 5.0

```
for (variable : collection)  
{  
    // loop body  
}
```

- **collection** is an array or an object of a class that implements **interface Iterable**
 - more on classes and interfaces later
- Executes the loop body for every element of the **collection**, assigned iteratively to **variable**

Control flow: **switch** selector



```
switch (Expr)
{
    case Value1: instructions;
        break;
    case Value2: instructions;
        break;
    // ...
    default: instructions;
}
```

Expr can be of type:

- **byte**, **short**, **int**, **char** (or wrapped counterparts)
- **enum** types
- **String** (compared with **equals**) (new in Java 7)

Breaking the control flow: **break**



label: [**while** | **do** | **for**]

- Identifies a loop
- (Or a code block)

break optionalLabel;

- Within a loop or a **switch**
- No label: exit the loop or switch
- With label:
 - within loop: jump out of the loop to label **optionalLabel**
 - within **switch**: jump out of **switch** block to label **optionalLabel**



`label: [while | do | for]`

- Identifies a loop
- (Or a code block)

`continue optionalLabel;`

- Within a loop
- No label: skip the remainder of the current iteration and continue with the next iteration
- With label:
 - skip the remainder of the current iteration and continue with the next iteration of the loop with label `optionalLabel`