





Part 1: Language constructs

1.3 EXCEPTION HANDLING

```
public class Printer {
   public print(int i) {
      try {
        throw new Exception()
      }
      catch(Exception e) { }
      }
```

Eiffel: Exception Handling

class PRINTER feature print_int (a_int: INTEGER) local I retried: BOOLEAN do if not / retried then (create {DEVELOPER EXCEPTI else -- Do something end rescue *I* retried := **True** retry end end OLD VIOLATION

EXCEPTION 🖃 🎱 ASSERTION VIOLATION CHECK_VIOLATION INVARIANT_VIOLATION LOOP INVARIANT VIOLATION POSTCONDITION VIOLATION PRECONDITION VIOLATION VARIANT_VIOLATION DEVELOPER EXCEPTION 🖃 🎱 MACHINE EXCEPTION 🖃 🎱 HARDWARE EXCEPTION FLOATING_POINT_FAILURE 🖃 🍉 OPERATING_SYSTEM_EXCEPTION COM_FAILURE OPERATING_SYSTEM_FAILURE OPERATING SYSTEM SIGNAL FAILURE OBSOLETE_EXCEPTION EXCEPTION_IN_SIGNAL_HANDLER_FAILURE RESCUE_FAILURE RESUMPTION FAILURE 🖃 🎱 SYS_EXCEPTION EIFFEL_RUNTIME_PANIC 🖃 🎱 EIF EXCEPTION 🖃 🎱 EIFFEL_RUNTIME_EXCEPTION 🖃 🍉 DATA_EXCEPTION IO FAILURE MISMATCH FAILURE SERIALIZATION_FAILURE EXTERNAL FAILURE NO_MORE_MEMORY 🖃 🎱 LANGUAGE EXCEPTION BAD INSPECT VALUE EIFFELSTUDIO_SPECIFIC_LANGUAGE_EXCEPTION ADDRESS APPLIED TO MELTED FEATURE CREATE_ON_DEFERRED ROUTINE FAILURE VOID_ASSIGNED_TO_EXPANDED VOID TARGET

```
feature
   transmit (a_p: PACKET)
        -- transmit packet a p
     local
        I current retries: INTEGER
        r: RANDOM_NUMBER_GENERATOR
      do
        line.send (a_p)
      rescue
        if I current retries < max retries then
           r.next
            wait_millisecs (r.value_between(20, 50))
            current retries := current retries + 1
           retry
        end
   end
end
```

 \mathbf{O}

Part 1: Language constructs

1.4 ONCE ROUTINES

```
foo: INTEGER

once

Result := factorial (10)

end

test_foo

do

io.put_integer (foo) -- 3628800, calculated

io.put_integer (foo) -- 3628800, directly returned

end
```

- Executed the first time
- Result is stored
- In further calls, stored result is returned
- In other languages
 - Static variables
 - Singleton pattern



Constants, other than basic types *i: COMPLEX*once create Result.make (0, 1) end

Lazy initialization settings: SETTINGS once create Result.load_from_filesystem end

>Initialization procedures
 init_graphics_system
 once ... end

Part 1: Language constructs

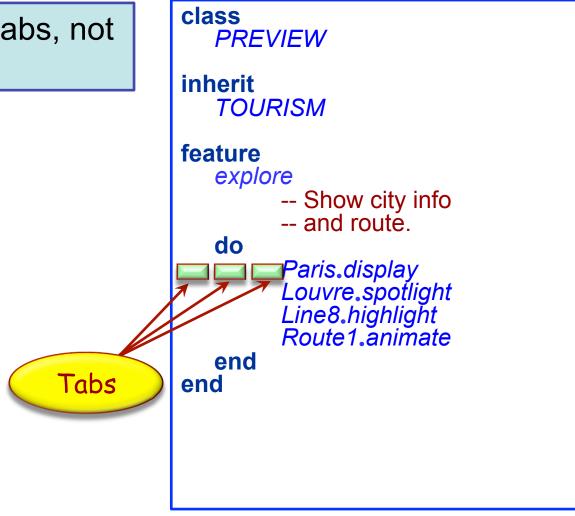




Style rule

For indentation, use tabs, not spaces





More style rules

- Class name: all upper-case Full words, no abbreviations (with some exceptions)

- Classes have global namespace: two classes cannot have the same name (even in different clusters)

- Usually, classes are prefixed with a library prefix

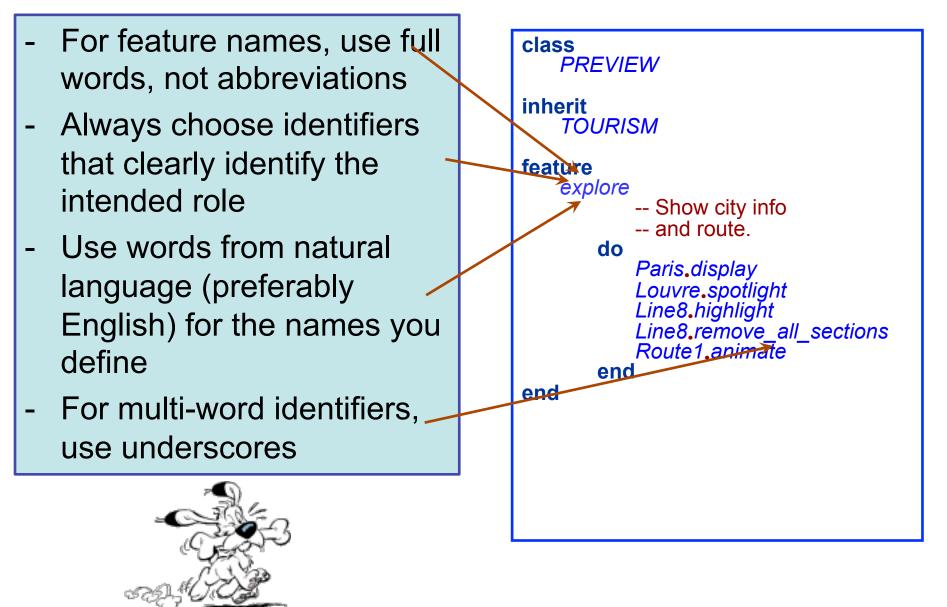
EiffelVision2: EV_

Base is not prefixed



class > PREVIEW inherit TOURISM feature explore -- Show city info -- and route. do Paris.display Louvre.spotlight Line8.highlight Route1.animate end end

Even more style rules



Eiffel Naming: Locals / Arguments

Locals and arguments share namespace with features

Name clashes arise when a feature is introduced, which has the same name as a local (even in parent)

➤To prevent name clashes:

- Locals are prefixed with
- Some exceptions like "i" exist
- Arguments are prefixed with a_



Part 1: Language constructs

1.6 GENERICS

class MY QUEUE [G]

feature

item: G

-- First item in queue.

do ... end

extend (a_element: G) -- Add new element. do ... end

end



G is called the generic parameter. By convention, the generic parameter name is G. If there are more parameters, use G, H, etc. or a meaningful abbreviation such as K for keys in a hash table

Creating instances of generics classes

class EXAMPLE1

feature

int_queue -- An integer queue. local qi: MY_QUEUE [INTEGER] do create qi qi.extend (35) qi.extend (6) end end

class EXAMPLE2

feature string_queue -- A string queue. local qs: MY_QUEUE [STRING] do create qs qs.extend ("Asterix") qs.extend ("Obelix") qs.extend ("Suffix") end end

class

MY_LIST [G -> COMPARABLE]

feature

item: G

-- First item in queue.

do ... end

. . .

```
extend (a_element: G)
-- Add new element.
do
```

```
... if a_element < item then
```

The generic parameter G must be a class inheriting from COMPARABLE



end

Creating instances of constraint generics classes

. . .

-- Valid declarations li: MY_LIST [INTEGER] ls: MY_LIST [STRING] lr: MY_LIST [REAL] ld: MY_LIST [DOUBLE]

. . .

-- Invalid declarations la: MY_LIST [ACCOUNT] lb: MY_LIST [BANK] lm: MY_LIST [MAIN]

> Classes ACCOUNT, BANK & MAIN don't inherit from COMPARABLE

Part 1: Language constructs

1.8 INFORMATION HIDING

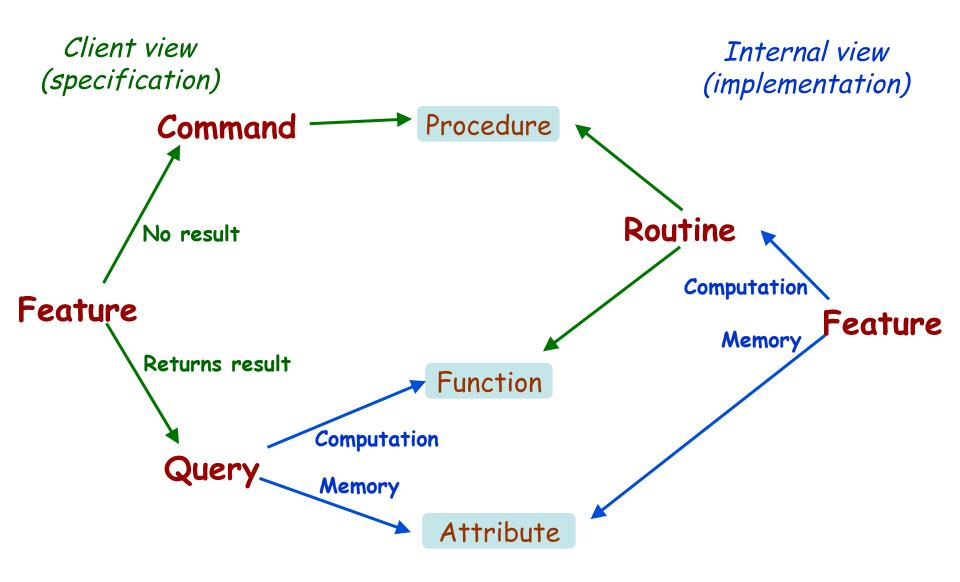


Procedure: doesn't return a result

- Yields a command
- Calls are instructions

Function: returns a result

- Yields a query
- Calls are expressions

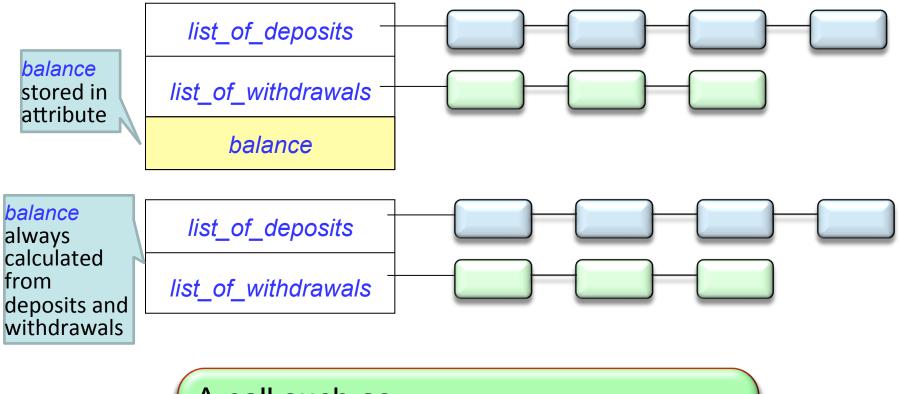


It doesn't matter to the client whether you look up or compute



Uniform Access: an example

balance = list_of_deposits.total - list_of_withdrawals.total



A call such as your_account.balance could use an attribute or a function

Exporting (making public) an attribute

In Eiffel, exporting an attribute means exporting it read-only

From the outside, it is not shown as an attribute, just as a **query**: it could be a function

In C++, Java & C#, if you make public an attribute* *x*, it is available for both read and write:

> a1.x := v

As a result, it is almost always a bad idea to export an attribute.

* (field, member variable)

 \mathbf{O}

Getter functions

In C++, Java & C#, the standard technique, if *private_x* is secret, is to export an associated **getter function**:

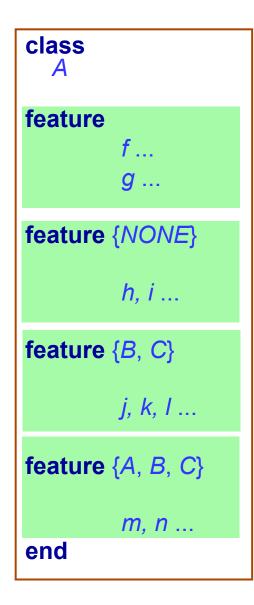
```
x : T
do
Result := private_x
end
```

Eiffel needs no getter functions: just export the attribute

This is safe: the attribute is exported

- Only for reading
- Without the information that it is an attribute: it could be a function (Uniform Access principle)

Information hiding



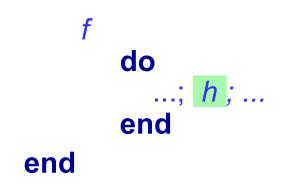
Status of calls in a client with a1: A:

- a1.f, a1.g: valid in any client
- a1.h: invalid everywhere
 - (including in A's own text!)
- a1.j: valid only in B, C and their descendants (not valid in A!)
- a1.m: valid in B, C and their descendants, as well as in A and its descendants

Information hiding only applies to use by clients, i.e. using dot notation or infix notation, as with *a1.f* (*Qualified* calls).

Unqualified calls (within class) not subject to information hiding:

```
class A feature {NONE }
h do ... end
feature
```



PART 2: CONTRACTS

Contracts

A contract is a semantic condition characterizing usage properties of a class or a feature

Three principal kinds:

- Precondition
- Postcondition
- Class invariant

Design by Contract

Together with the implementation ("*how*") of each software element, describe "*what*" it is supposed to do: its contract

Three basic questions about every software element:

What does it assume?

> What does it guarantee?
> What does it maintain?
Precondition
Postcondition

Contracts in programming languages

Eiffel: integrated in the language

Java: Java Modeling Language (JML), iContract etc.

.Net languages: Code Contracts (a library)

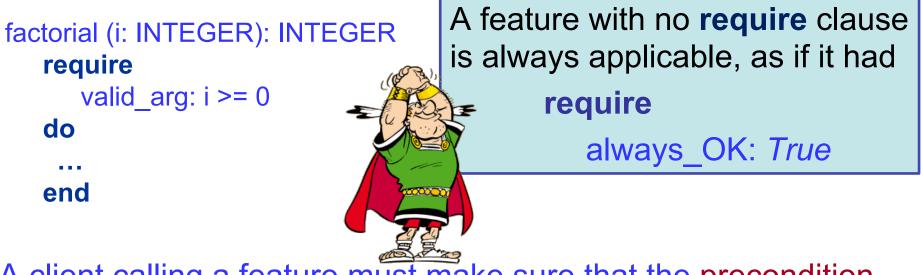
Spec# (Microsoft Research extension of C#): integrated in the language

UML: Object Constraint Language

etc.

Precondition

Property that a feature imposes on every client:



A client calling a feature must make sure that the precondition holds before the call

A client that calls a feature without satisfying its precondition is faulty (buggy) software.

Another example:

```
extend (a_element: G)

require

valid_elem: a_element /= void

not_full: not is_full

do ... end
```

```
A feature with a require clause

require

label_1: cond_1

label_2: cond_2 ...

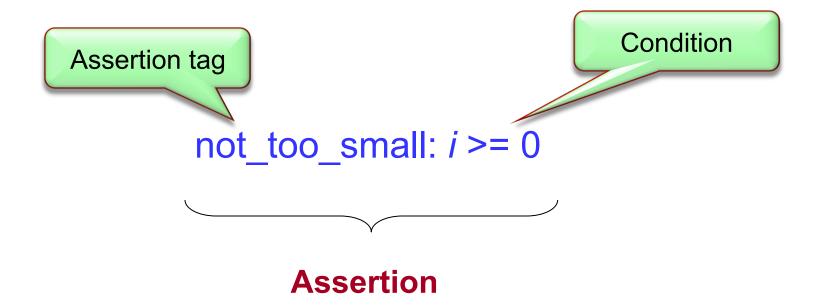
label_n: cond_n

is equivalent to

require

label: cond_1 and cond_2 and ... cond_n
```





 \odot

Let's code...

Go to:



https://codeboard.io/projects/86

Task: in class CUSTOMER, write a precondition for the creation routine make_with_name_and_age

Task: create an invalid CUSTOMER object and try to run your program. What happens?

Task: fix your **CUSTOMER** object to satisfy the precondition of the creation routine

Postconditions

Precondition: obligation for clients Postcondition: benefit for clients

```
extend (a_element: G)
ensure
inserted: i_th (count) = a_element
```

```
index (a_element: G): INTEGER
```

ensure

```
exists: result > 0 implies i_th (result) = a_element
no_exists: result = -1 implies not is_inserted (a_element)
```

 \mathbf{O}

Let's code...

Go to:



()

https://codeboard.io/projects/86

Task: in class CUSTOMER, write a postcondition for the creation routine make_with_name_and_age

Task: modify the implementation of make_with_name_and_age such that it breaks your postcondition. Run the program. What happens?

Old notation

Usable in postconditions only

Denotes value of an expression as it was on routine entry

Example (in a class ACCOUNT):

```
balance: INTEGER
              -- Current balance.
deposit (v: INTEGER)
              -- Add v to account.
       require
              positive: v > 0
       do
       ensure
              added: balance = old balance + v
       end
```

A feature must make sure that, if its precondition held at the beginning of its execution, its postcondition will hold at the end.

A feature that fails to ensure its postcondition is buggy software.



Invariant

An invariant states properties about an object that are true

- after the object has been initialized
- before and after every routine call (but not necessarily in between a call)

The invariant is listed after the last feature block.

A class with no **invariant** is the same a **invariant** always_OK: *True*

Example (from class ARRAY):

invariant

area_exists: area /= Void consistent_size: capacity = upper - lower + 1 non_negative_count: count >= 0 index_set_has_same_count: valid_index_set

class

BANK_ACCOUNT create make feature make (n : STRING) -- Set up with name n

require n /= Void

do

name := n balance := 0

ensure name = n

end

name : STRING balance : INTEGER deposit (v: INTEGER) -- Add amount v do balance := balance + vensure balance = old balance + vend invariant name /= Void

balance >= 0

end

Let's code...

Go to:



()

https://codeboard.io/projects/86

Task: in class ACCOUNT, replace all -- Important: ... comments with contracts

Contracts and inheritance (Example)

class	ACCOUNT
ACCOUNT_MANAGER	ACCOUNT_
feature Operations	MANAGER
<pre>init_new_account(a_acc: ACCOUNT)</pre>	
do all initialization of po	not strengthen ondition because lymorphism and mic binding.
	class
class ACCOUNT	SPECIAL_ACCOUNT inherit ACCOUNT redefine set_balance end
feature Operations	feature Operations
<pre>set_balance(a_balance: DOUBLE) require non_neg: a_balance >= 0 do balance := a_balance</pre>	<pre>set_balance(a_balance: DOUBLE) require min_bal: a_balance > 100 do</pre>
end	<pre>balance := a_balance end</pre>

When redeclaring a routine, we may only:

Keep or weaken the precondition

Keep or strengthen the postcondition

Invariant Inheritance rule:

The invariant of a class automatically includes the invariant clauses from all its parents,

"and"-ed.

A simple language rule does the trick!

Redefined version may have nothing (assertions kept by default), or

require else new_pre
ensure then new_post

Resulting assertions are:

original_precondition or new_pre

original_postcondition and new_post

 \mathbf{O}