# Bita: Coverage-Guided, Automatic Testing of Actor Programs

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### Background

- 2 Implementation
- 3 Evaluation
- 4 Conclusions

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# **Testing Concurrent Programs**

- Potentially a large number of different interleavings of operations.
- Test may succeed with some interleavings and fail with others.
- Requirements:
  - Test different interleavings.
  - Testing should not take too long.

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# Actors (in Akka)

- Program is a set of actors.
- Actor entity with its own local state and thread of control that communicate exclusively by exchanging messages.
  - A mail box for incoming messages.
  - A message handler, which can *change at runtime*.
  - Message processing is *atomic*.
- **Execution schedule** order in which actors receive messages.

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### Overview

- **1** Select coverage criterion.
- 2 Obtain initial schedule by running the program with default scheduler.
- **3** Generate interesting schedules by reordering initial schedule.
  - "Interesting" = increases coverage.
  - Goal: only feasible schedules.
  - Goal: minimize the amount of schedules generated.
- 4 Run program with generated schedules.

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# Coverage Criteria

- A pair of receive events for the *same* receiver.
  - 1 Pair of Consecutive Receives (PCR).
  - 2 Pair of Receives (PR).
  - 3 Pair of Message Handler Change and Receive (PMR).
- For the specific coverage criterion a set of schedules covers a pair of receive events if and only if there exists schedules that cover both orderings.

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# Must-Happen-Before Constraints

- 1 Causality Constraints one event caused another.
- Sender-Receiver Constraints messages between two actors are delivered in order.
- **3** Ordering Constraints synchronous communication.

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# Schedule Generation Algorithm (Simplified)

1:	function Schedul	E(prefix, tail, cr)		
2:	for all $r_i, r_j \in ta$	$il \wedge r_i$ before $r_j$ <b>do</b>		
3:	if isCrRelate	$d(r_i, r_j, cr) \land (r_i, r_j) \not\in m$	nustHB <b>then</b>	
4:	if $r_i \rightarrow_{cr}$	$r_j \notin OrderingGoals$ the	n	
5:	newP	$Prefix \leftarrow prefix + before($	$(r_i) + mustHB(r_j) + r_i -$	⊦ r <sub>j</sub>
6:	newT	$ail \leftarrow \dots$		
7:	Ordei	$ringGoals \leftarrow OrderingGoals$	$pals \cup \{r_i \to_{cr} r_j\}$	
8:	retur	<b>n</b> Schedule( <i>newPrefix</i>	r, newTail, cr)	
9:	end if			
10:	if $r_j \rightarrow_{cr}$	$r_i \notin OrderingGoals$ the	n	
11:	end if			
12:	end if			
13:	end for			
14:	<b>return</b> prefix			
15:	end function			
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# **Bug Detection**

	Issue	Bita		Random Scheduler		Default Scheduler		
Bug					$d_{max}=300$ ms			
e		Tried Criteria	Time	Schedule	Time	Execs	Time	Execs
Ga1(U)	1019	PR	36±1	1	ТО	191	ТО	265
Ga2(U)	1018	PR	$37 \pm 1$	1	$163 \pm 75$	$8\pm4$	ТО	269
Ga3(U)	1018	PR	$26 \pm 1$	1	$100 \pm 40$	$6\pm 2$	TO	270
Ga4(U)	1116	PR	$25 \pm 1$	1	$326 \pm 98$	$18 \pm 5$	ТО	270
SC1(U)	80	PR	$102 \pm 15$	$2\pm 1$	TO	158	TO	182
SC2(U)	81	PR	$86 \pm 32$	$2\pm 1$	ТО	104	ТО	219
SC3(K)	58	PR	$176 \pm 29$	$3\pm 1$	ТО	90	ТО	257
FR11(U)	13	PR	$43 \pm 6$	1	TO	206	ТО	225
FR12(K)	12	PR	$36 \pm 1$	1	TO	471	ТО	594
Ba(U)		PR,PMR	$250 \pm 43$	$28\pm5$	TO	263	ТО	532
Ms(K)		PR	14	1	TO	703	TO	1788
PR(K)		PR,PCR	$263 {\pm} 151$	$32\pm21$	TO	235	$2,268{\pm}782$	$557 \pm 180$
		Summary of	all buge wi	th ton ron	atitions nor l	bua:		
		Summary of	all bugs wi	-	-	bugs—Avg. time	to detect a bug	g—Slowdown
10,939—120—91— <b>1</b> x   335,903—30—11,196— <b>122</b> x   419,020—7—								9,860 <b>—656x</b>

Table II. Times are in seconds. "TO" - timeout (1 hour).

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### Criteria Comparison

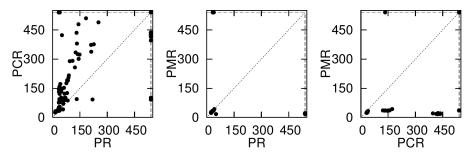


Fig. 3. Pairwise comparison of time (in seconds) needed to detect a bug with specific criterion.

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### Coverage

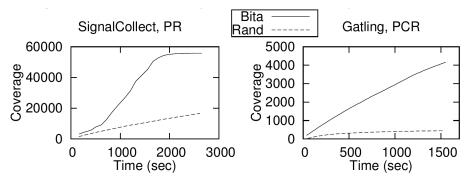


Fig. 4. Comparison of coverage achieved by Bita and random scheduling with  $d_{max} = 300 ms$ .

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# Conclusions

### Advantages:

- Bita is much faster in finding bugs than alternatives.
- Schedules that reveal bugs are logged.
- Limitations:
  - Schedules are generated based on the single run data.
  - Conservative must-happen-before constraints.
- Impact:
  - Paper is cited in 3 papers, but only in related or future work sections.
- Unclear points:
  - What setup (hardware, OS, JVM) was used for experiments?
  - Is final test execution parallel?

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Conclusions

# Thank You!

Questions?

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