

# Problem Sheet 6: Software Model Checking

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The exercises in this problem sheet are based on the software model checking slides:

[http://se.inf.ethz.ch/courses/2015b\\_fall/sv/slides/11-SoftwareModelChecking.pdf](http://se.inf.ethz.ch/courses/2015b_fall/sv/slides/11-SoftwareModelChecking.pdf)

## 1 Predicate Abstraction

Recall that:

- $Pred(f)$  denotes the weakest under-approximation of the expression  $f$  expressible as a Boolean combination of the given predicates.
- The Boolean abstraction of an `assume  $c$  end` statement is `assume not  $Pred(not\ c)$  end` followed by a parallel conditional assignment updating the predicates with respect to the original `assume` statement.
- The Boolean abstraction of an `assert  $c$  end` statement is `assert  $Pred(c)$  end`.

### Exercises

- Justify, using Venn diagrams, the use of double negation in the Boolean abstraction of `assume` statements.
- Consider the following code snippet  $C_1$ , where  $x, y, z$  are integer variables:

```
assume x > 0 end
z := (x*y) + 1
assert z >= 1 end
```

Build the Boolean abstraction  $A_1$  of the code snippet  $C_1$  with respect to the following set of predicates:

$$p \triangleq x > 0$$

$$q \triangleq y > 0$$

$$r \triangleq z > 0$$

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\*Some exercises were adapted from earlier ones written by Stephan van Staden and Carlo A. Furia.

- iii. Consider the following code snippet  $C_2$ , where  $x, y$  are integer variables:

```
if x > 0 then
  y := x + x
else
  if x = 0 then
    y := 1
  else
    y := x * x
  end
end
assert y > 0 end
```

- (a) Normalise the guards of conditionals using nondeterminism and **assume** statements.  
(b) Build the Boolean abstraction  $A_2$  of the normalised code snippet  $C_2$  with respect to the following set of predicates:

$$p \triangleq x > 0$$

$$q \triangleq y > 0$$

## 2 Error Traces

- i. Provide an annotated trace for the Boolean abstraction  $A_1$ , and a corresponding (feasible) annotated trace for the concrete program  $C_1$  in which **assert z >= 1 end** evaluates to false when reached.
- ii. Can you verify the Boolean abstraction  $A_2$ ? If not, give a trace as a counterexample and prove whether or not it is *spurious*.