

Software Verification (Autumn 2015)

Lecture 5: Auto-Active Verification

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(based on material by Nadia Polikarpova)

This time last week

$\vdash \{x > 0\} x := x + 1; \text{skip} \{x > 1\}$

$$\begin{array}{c} \text{[ass]} \frac{}{\vdash \{x + 1 > 1\} \mathbf{x} := \mathbf{x} + \mathbf{1} \{x > 1\}} \\ \text{[cons]} \frac{}{\vdash \{x > 0\} \mathbf{x} := \mathbf{x} + \mathbf{1} \{x > 1\}} \\ \text{[comp]} \frac{}{\vdash \{x > 0\} \mathbf{x} := \mathbf{x} + \mathbf{1}; \text{skip} \{x > 1\}} \end{array} \quad \begin{array}{c} \text{[skip]} \frac{}{\vdash \{x > 1\} \text{skip} \{x > 1\}} \end{array}$$

Can we reason about $\{\text{pre}\}P\{\text{post}\}$ mechanically?

Verification problem **undecidable in general**

How far can we go? What challenges do we face?

- Determining loop invariants

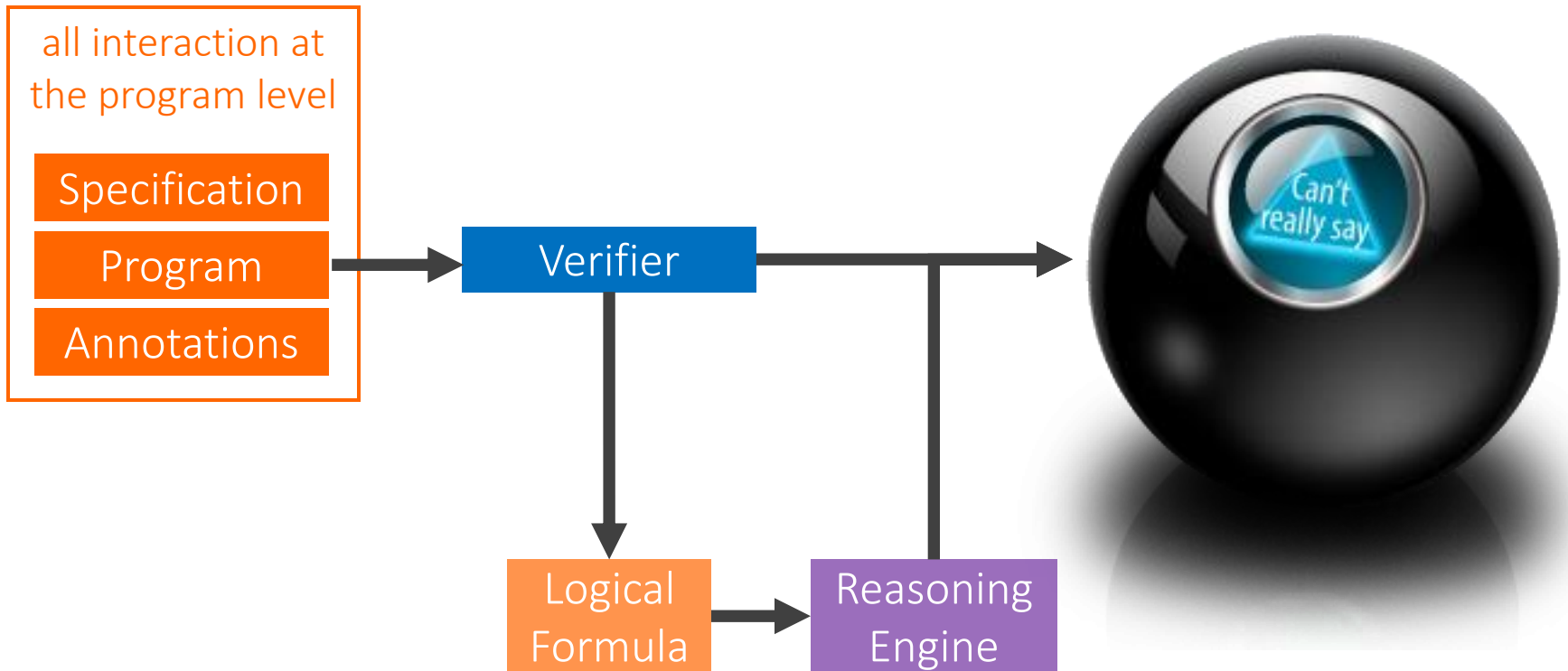
- Weak or missing assertions

- Undecidable assertion logics

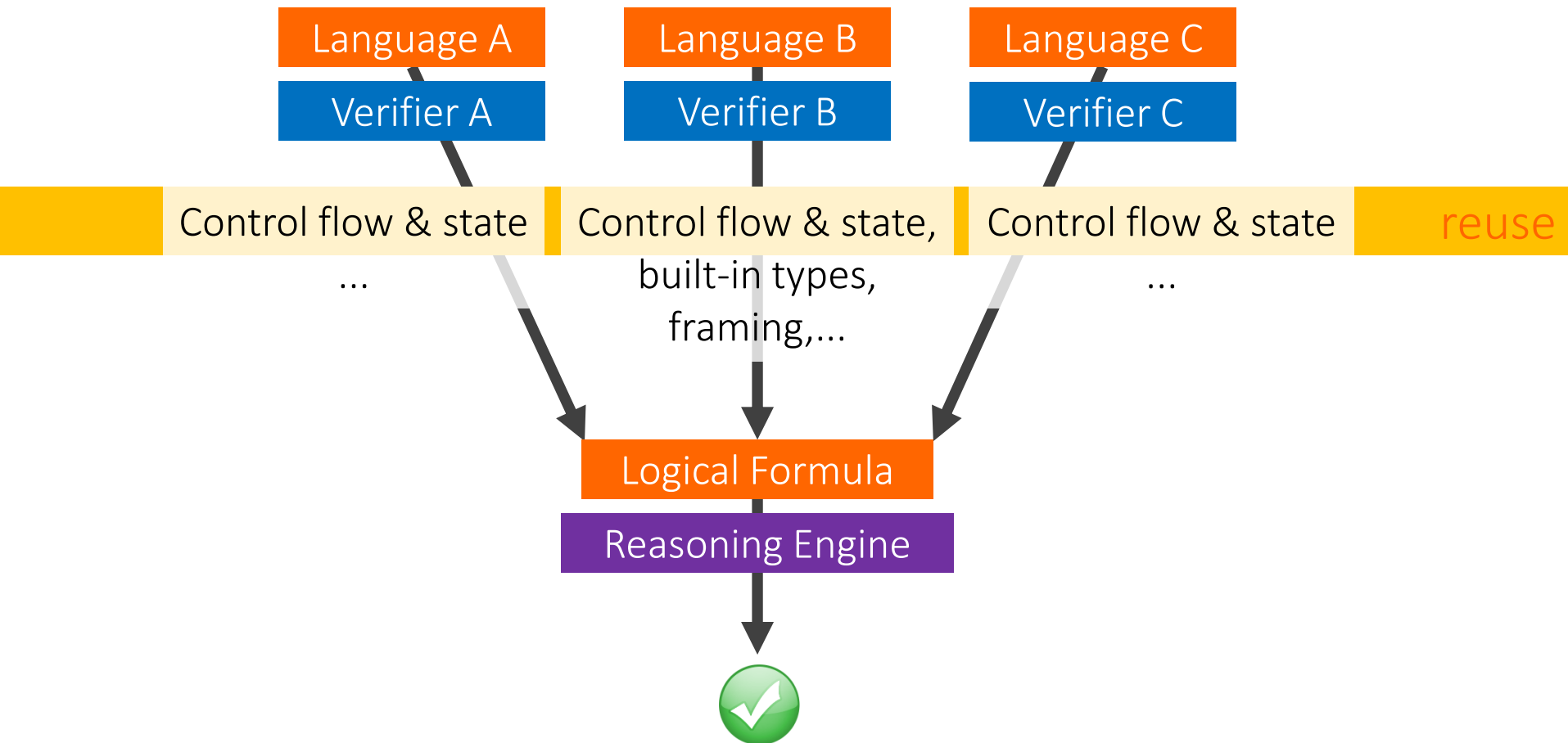
- ...

Idea: automate as much as possible, with users indirectly providing guidance through program-level annotations

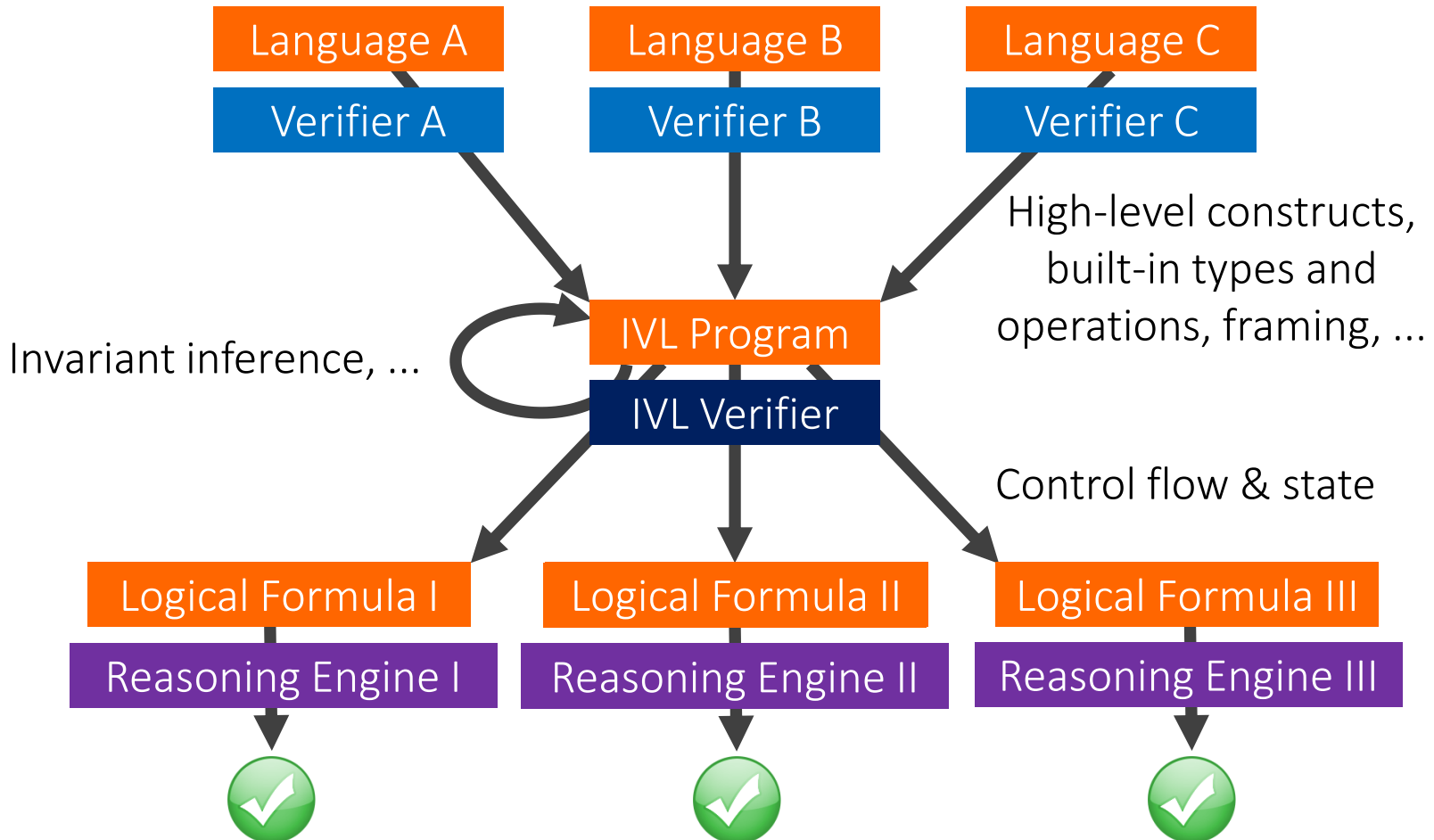
“Auto-active” verification



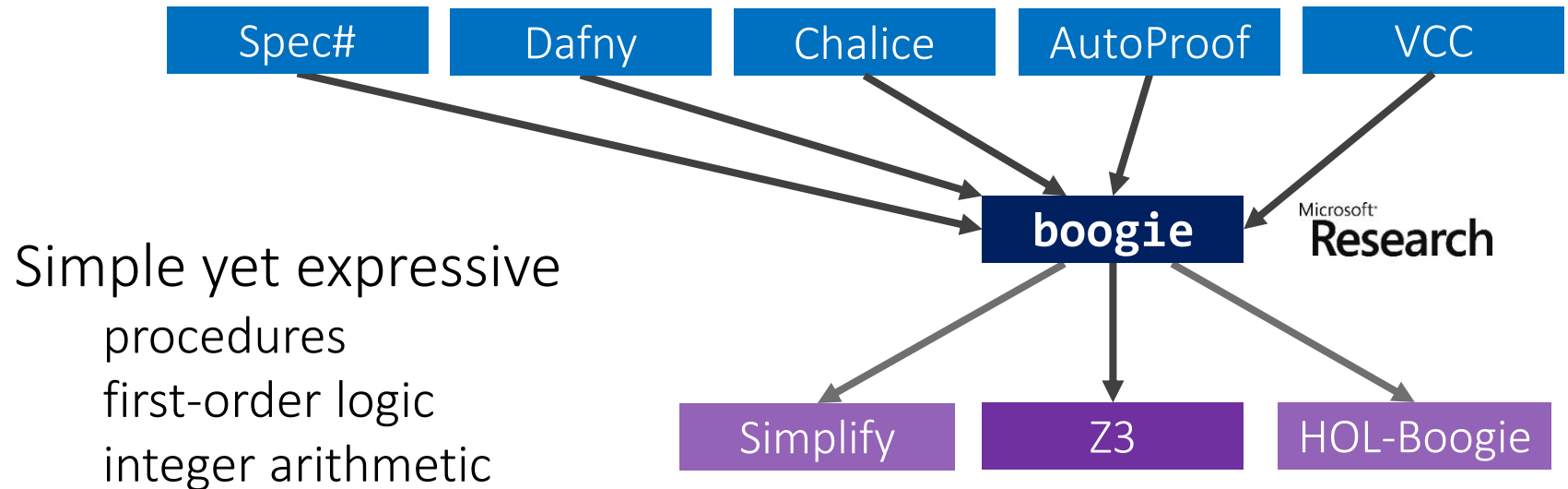
Verifying imperative programs



Intermediate Verification Language



The Boogie IVL



Great for teaching verification!

skills transferable to other auto-active tools

Alternatives: Why3 [<http://why3.lri.fr/>]

Viper [<http://www.pm.inf.ethz.ch/>]

Overview

The Boogie Language : **how to express your intention?**

- Imperative constructs

- Specification constructs

The Boogie Tool : **how to get it to verify?**

- Debugging techniques

- Boogaloo to the rescue

The AutoProof Verifier

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Getting started with Boogie

boogie Microsoft
Research

Try online [rise4fun.com/Boogie]

Download [boogie.codeplex.com]

User manual [Leino: [This is Boogie 2](#)]

[Hello, world?](#)

Types

Basic types: **bool**, **int**, **real**

definition

User-defined: **type** Name t_1, \dots, t_n ;

usage

type ref; // references

type Person;

type Field t; // fields with values of type t

Field **int**

Field ref

Maps: $\langle t_1, \dots, t_n \rangle [dom_1, \dots, dom_n] range$

[**int**]int // array of int

[Person]**bool** // set of persons

[ref]ref // “next” field of a linked list

$\langle t \rangle [ref, Field t]t$ // generic heap

Synonyms: **type** Name $t_1, \dots, t_n = type$;

type Array t = [int]t;

type HeapType = $\langle t \rangle [ref, Field t]t$;

Imperative constructs

Regular procedural programming language

[\[Absolute Value & Fibonacci\]](#)

... and non-determinism

great to simplify and over-approximate behavior

```
havoc x; // assign an arbitrary value to x
```

```
if (*) { // choose one of the branches non-deterministically  
    statements  
} else {  
    statements  
}
```

```
while (*) { // loop some number of iterations  
    statements  
}
```

Specification statements: **assert**

assert *e*: executions in which *e* evaluates to **false** at this point are **bad**

expressions in Boogie are pure, no procedure calls

Uses

explaining semantics of other specification constructs

encoding requirements embedded in the source language

```
assert lo <= i && i < hi; // bounds check  
result := array[i];
```

```
assert this != null; // 0-0 void target check  
call M(this);
```

debugging verification (see later)

[\[Absolute Value\]](#)

Specification statements: `assume`

`assume` `e`: executions in which `e` evaluates to **`false`** at this point are **impossible**

```
havoc x; assume x*x == 169; // assign such that
```

```
assume true; // skip
```

```
assume false; // this branch is dead
```

Uses

- explaining semantics of other specification constructs
- encoding properties guaranteed by the source language

```
havoc Heap; assume NoDangling(Heap); // managed language
```

- debugging verification (see later)

Assumptions are dangerous! [[Absolute Value](#)]

Loop invariants

```
before_statements;  
while (c)  
    invariant inv;  
{  
    body;  
}  
after_statements;
```

=

```
before_statements;  
assert inv;  
  
havoc all_vars;  
assume inv && c;  
body;  
assert inv;  
  
havoc all_vars;  
assume inv && !c;  
after_statements;
```

The only thing the verifier know about a loop
simple invariants can be inferred

[[Fibonacci](#)]

Procedure contracts

```
procedure P(ins) returns (outs)  
  free requires pre';  
  requires pre;  
  modifies vars; // global  
  ensures post;  
  free ensures post';  
{ body; }
```

```
call outs := P (ins);
```

=

```
assume pre && pre';  
body;  
assert post;
```

=

```
assert pre;  
havoc outs, vars;  
assume post && post';
```

The only thing the verifier knows about a call

this is called **modular verification**

[[Abs and Fibonacci](#)]

Enhancing specifications

How do we express more complex specifications?

e.g. `ComputeFib` actually computes Fibonacci numbers

Uninterpreted functions

```
function fib(n: int): int;
```

Define their meaning using axioms

```
axiom fib(0) == 0 && fib(1) == 1;  
axiom (forall n: int :: n >= 2 ==> fib(n) == fib(n-2) + fib(n-1));
```

[[Fibonacci](#)]

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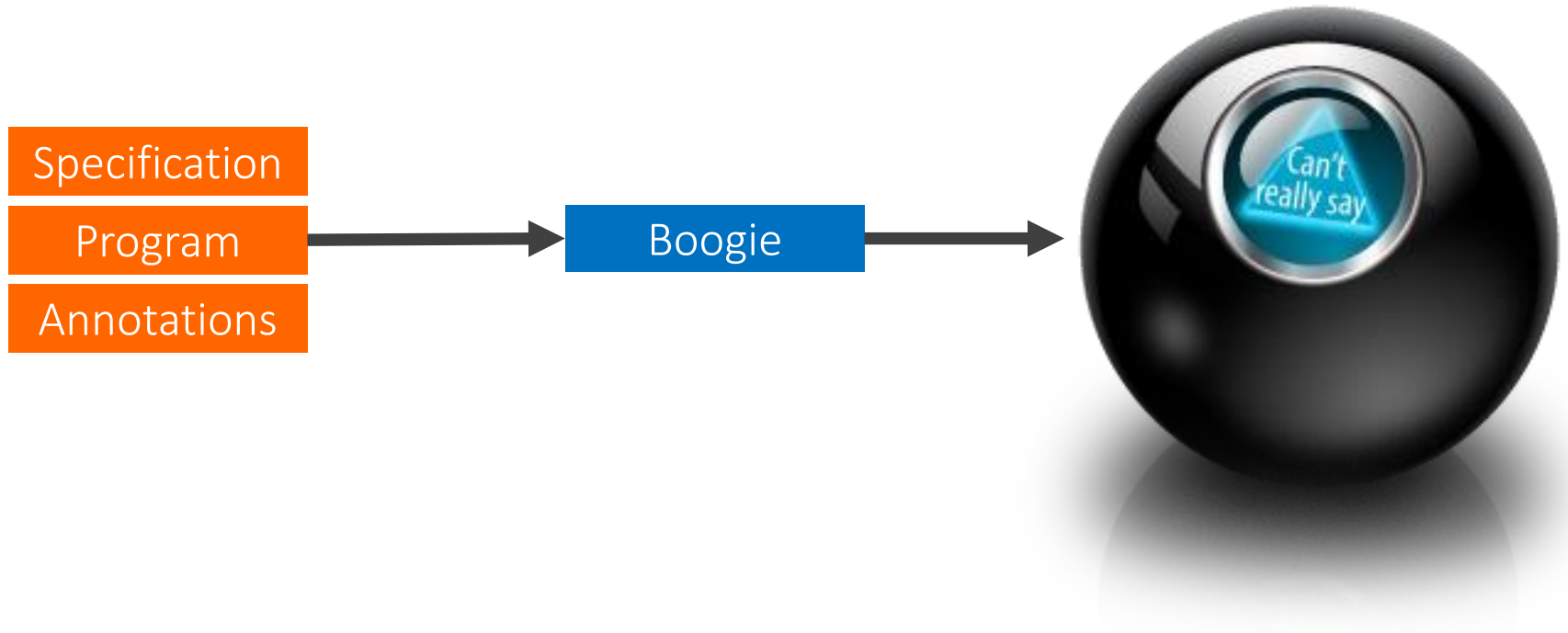
The Boogie Tool

- Debugging techniques

- Boogaloo to the rescue

The AutoProof Verifier

What went wrong?



Debugging techniques

Proceed in small steps [[Swap](#)]

use **assert** statements to figure out what Boogie knows

Divide and conquer the paths

use **assume** statements to focus on a subset of executions

Prove a lemma [[Non-negative Fibonacci](#)]

write ghost code to help Boogie reason

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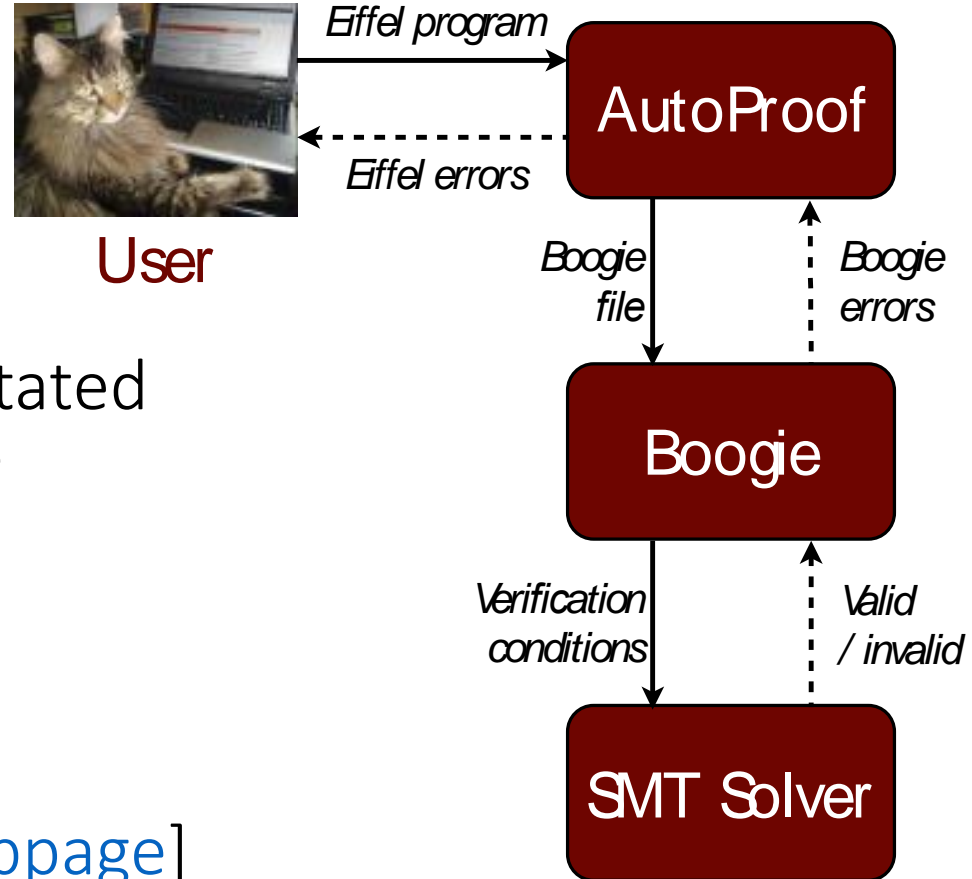
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The AutoProof Verifier

AutoProof: a Boogie-based verifier for Eiffel



Translates contract-annotated Eiffel programs to Boogie

Try online [[via Comcom](#)]

Manual, tutorial,

examples [[AutoProof webpage](#)]

How the translation works [[Slides](#)]

```

1  |-- Simple bank account class.
2  -- Try to fix it and make the verification go through.
3
4  class
5      ACCOUNT
6
7  feature -- Access
8
9      balance: INTEGER
10     | | | -- Balance of account.
11
12 feature -- Element change
13
14     deposit (amount: INTEGER)
15     | | | -- Deposit `amount` on account.
16     | | | require
17     | | |     amount_not_negative: amount >= 0
18     | | | do
19     | | |     balance := balance + amount
20     | | | ensure

```

Run...

Feature	Line	Result
ACCOUNT (invariant admissibility)		Successfully verified.
ANY.default_create (creator, inherited by ACCOUNT)		Successfully verified.
ACCOUNT.deposit		Successfully verified.
ACCOUNT.withdraw	31	Postcondition balance_decreased may be violated.
ACCOUNT.transfer	49	Postcondition balance_decreased may be violated.

Conclusions

Boogie is an **Intermediate Verification Language** (IVL)

IVLs help develop verifiers

The Boogie **language** consists of:

imperative constructs \approx Pascal

specification constructs (**assert**, **assume**, **requires**, **ensures**, **invariant**)

math-like part (functions + first-order axioms)

There are several **techniques** to debug a failed verification attempt

AutoProof is one of several **auto-active verifiers**, based on translating annotated programs to Boogie